

Constraint Based Synthesis

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Synthesis: 1980s view

Complete Formal Specification

Theory *Sorting*($\langle \alpha, \leq \rangle$: linear - order)
Imports *integer*, *bag*(α), *seq*(α)
Operations
 Ordered : *seq*(α) \rightarrow Boolean
Axioms
 $\forall(S : \text{seq}(\alpha)) (\text{Ordered}(S) \Leftrightarrow \forall(i)(i \in \{1..length(S) - 1\} \Rightarrow S(i) \leq S(i + 1)))$



```
int[N] sort(int N, int[N] input){
  int[N] output=input;
  int[N] done = 0;
  int k=0;
  for(int i=0; i<N; ++i){
    for(int j=i+1; j<N; ++j){
      if( output[j]< output[i]){
        int tmp = output[j];
        output[j] = output[i];
        output[i] = tmp;
      }
    }
  }
  return output;
}
```

Synthesis: modern view

Space of programs

 P_c $\varphi(c)$

Safety Properties

Input/Output Examples

Test Harnesses

Sketch

C-like language
with holes and
assertions

 P_c

```
/* Average of x and y without overflow */  
int avg(int x, int y){  
    int t =  
  
    return t;  
}
```

```
/* Average of x and y without overflow */
int avg(int x, int y){
    int t = y / 2 + x / 2 + ((x % 2 + y % 2) / 2);
    assert t == (x+y)/2;
    return t;
}
```



P_c

Automated
Tutoring



P_c

Automated
Tutoring

Provably
Correct
Synthesis

Program
Optimization



P_c

Automated
Tutoring

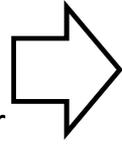
Provably
Correct
Synthesis

Program
Optimization

Visual
Programming

Program
+
Properties

□ Abstract
Domain
Compiler



Sketch
C-like language
with holes and
assertions

□ Unroll
Inline
Enumerate



$\varphi(c)$

P_c

Automated
Tutoring

Provably
Correct
Synthesis

Program
Optimization

Visual
Programming

Automated Tutoring

With Rishabh Singh and Sumit Gulwani



Pinpoint student errors

```
def computeDeriv(poly):  
    deriv = []  
    zero = 0  
    if (len(poly) == 1):  
        return deriv  
    for e in range(0, len(poly)):  
        if (poly[e] == 0):  
            zero += 1  
        else:  
            deriv.append(poly[e]*e)  
    return deriv
```

replace derive by [0]

Should be a 1

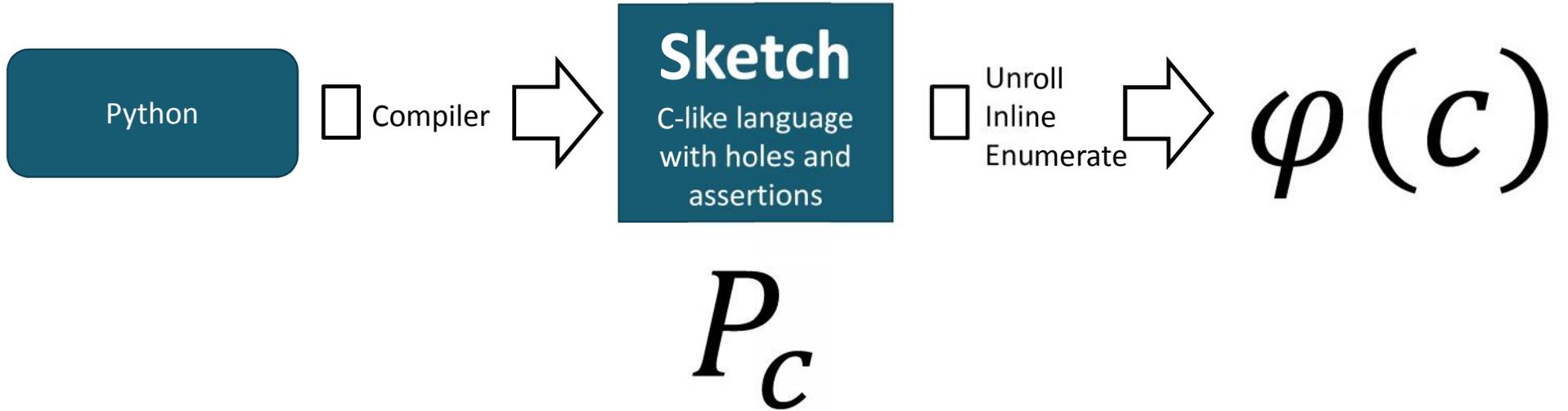
Get rid of this

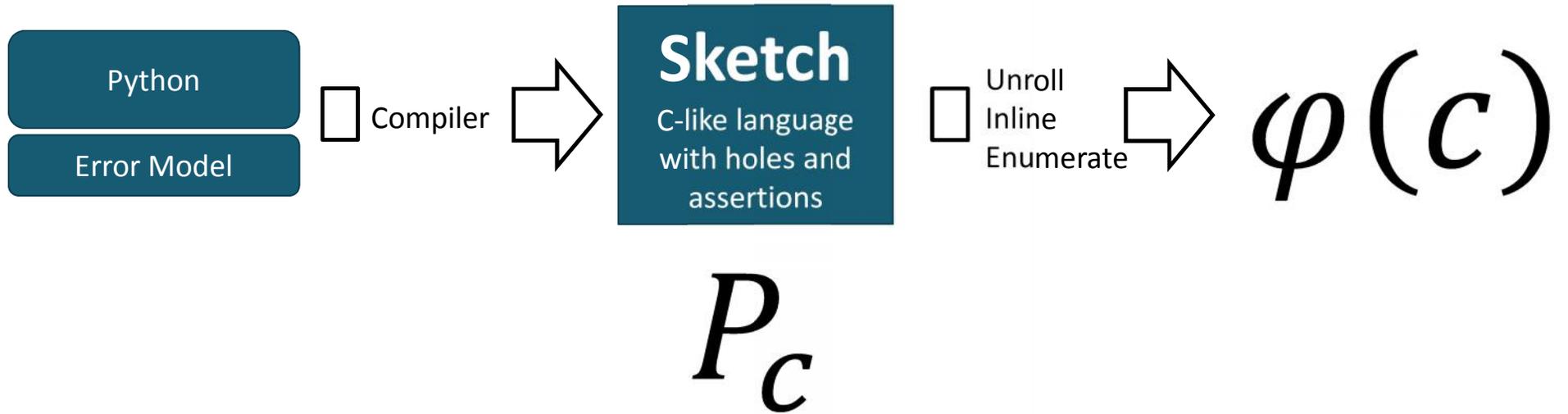
```
def computeDeriv(poly):  
    length = int(len(poly)-1)  
    i = length  
    deriv = range(0, length)  
  
    if len(poly) == 1:  
        deriv = [0.0]  
    else:  
        while i >= 0:  
            new = round((poly[i] * i), 2)  
            i -= 1  
            deriv[i] = new  
  
    return deriv
```

Should be a 0

Should be >

This is a synthesis problem





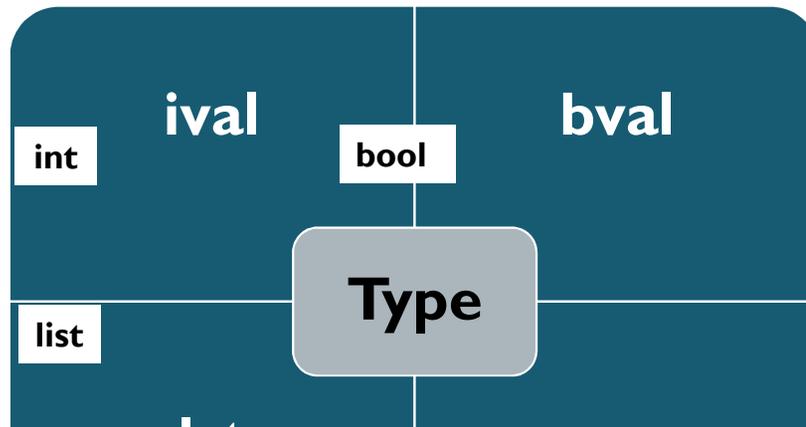


Dynamic Types

```

struct MultiType{
    int type;
    int ival;
    bit bval;
    MTString str; MTList lst;
    MTDict dict; MTTuple tup;
}

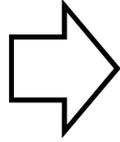
```



Python



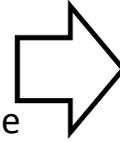
Compiler



Sketch

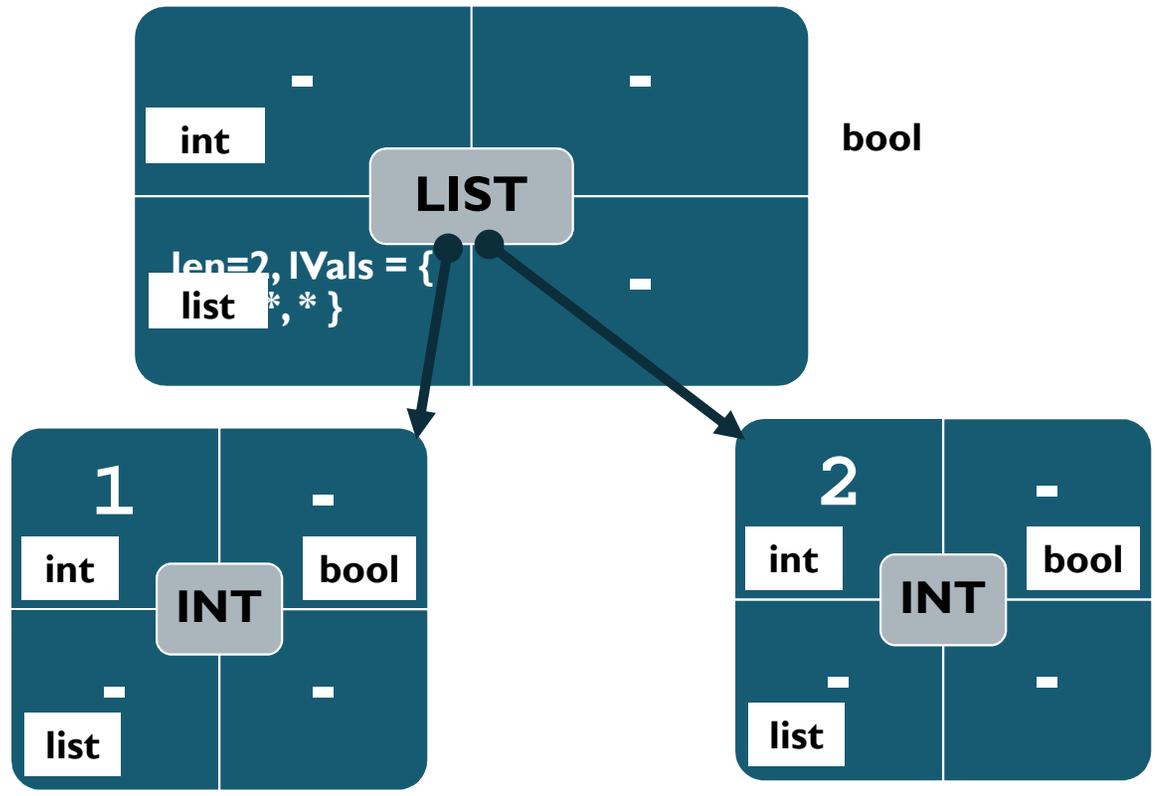


Unroll
Inline
Enumerate



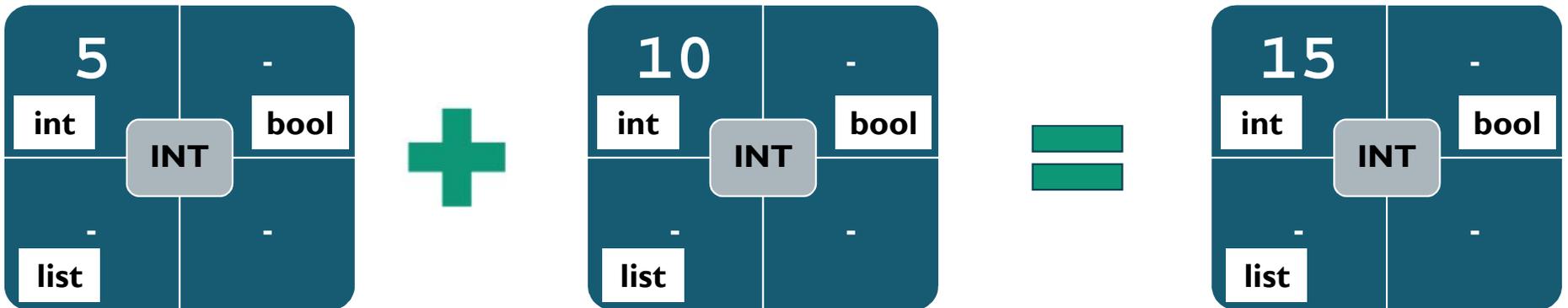
$\varphi(c)$

[1, 2]



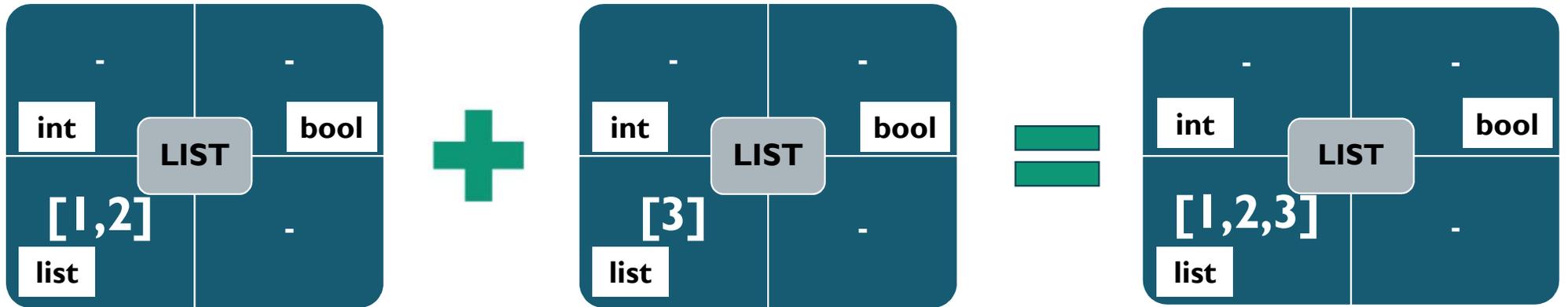


$x + y \rightarrow \text{addMT}(x, y)$



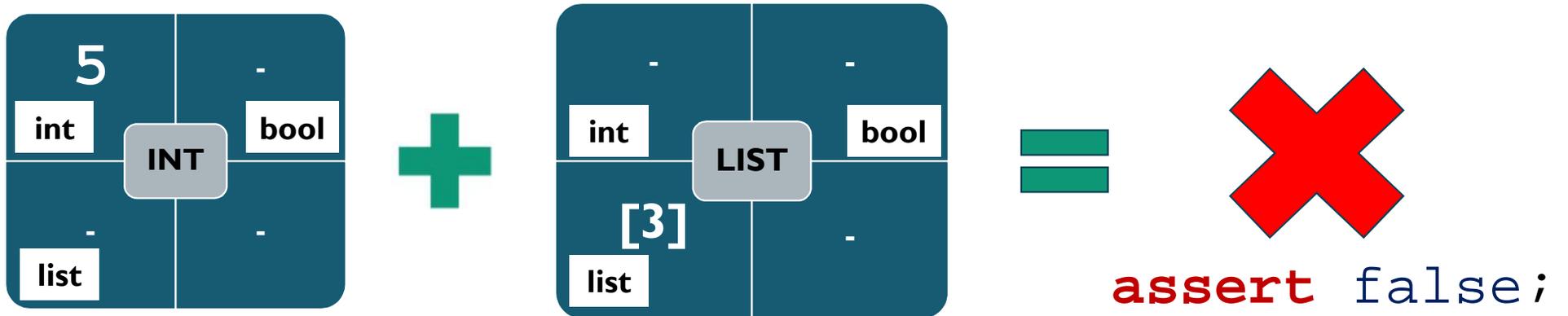


$\mathbf{x} + \mathbf{y} \rightarrow \text{addMT}(\mathbf{x}, \mathbf{y})$





`x + y` \rightarrow `addMT(x, y)`





- `return a` \rightarrow `return [0]`
- `range(a1, a2)` \rightarrow `range(a1+1, a2)`
- `a0 == a1` \rightarrow `False`



```
range(0, len(poly))
```



`range(a1, a2) → range(a1+1, a2)`

```
range({0}, 1, len(poly))
```

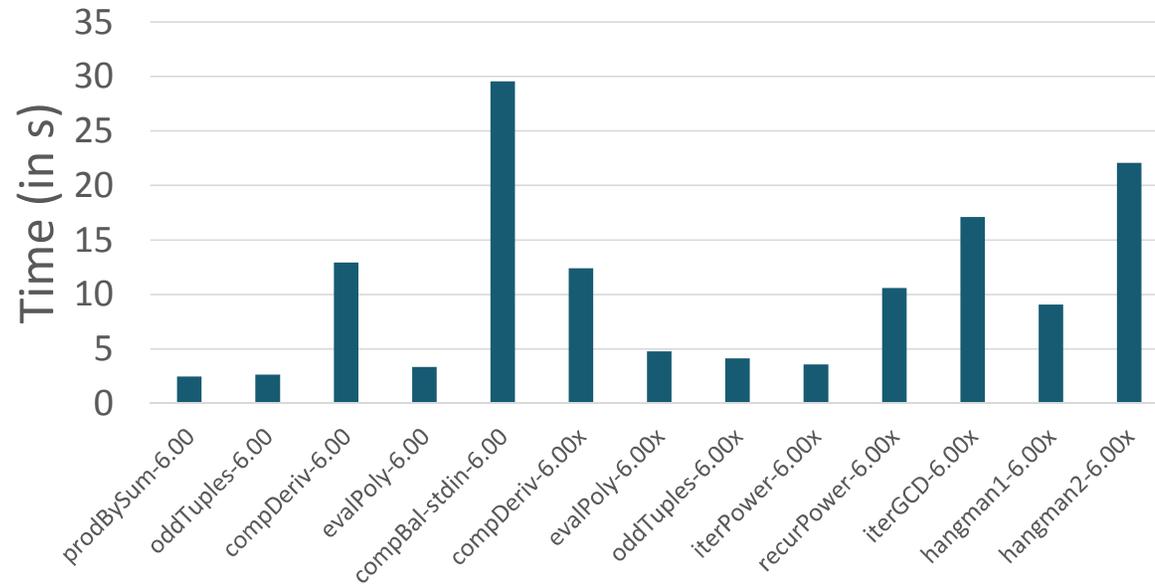
default choice



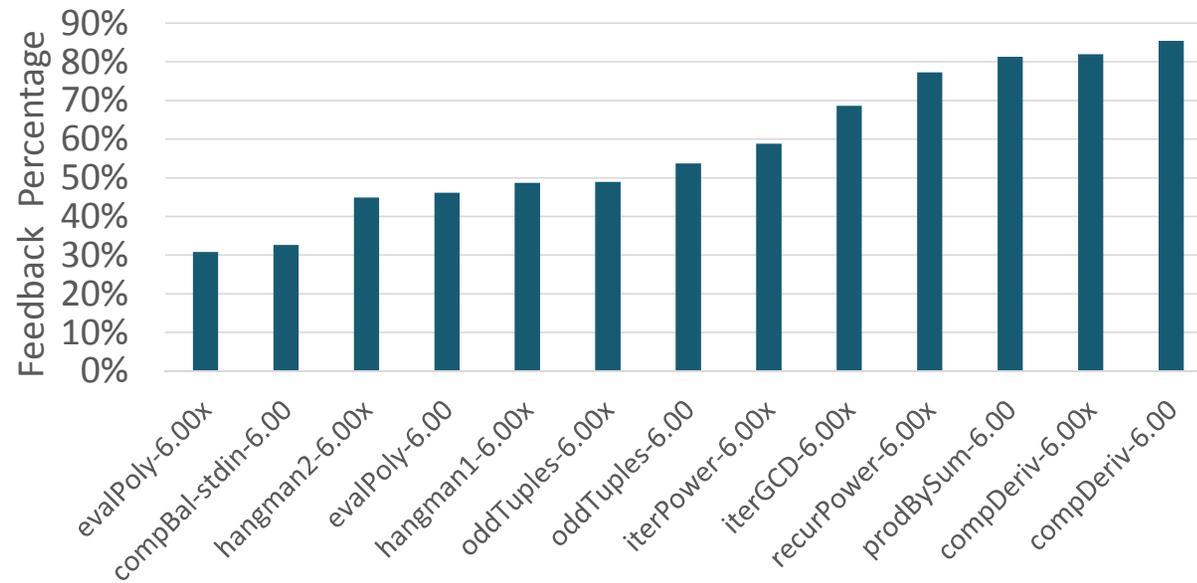
 ,  } \rightarrow modifyMT0()

```
MultiType modifyMT0() {
    if(??) return  // default choice
    else
        choice0 = true
        return  // non-default choice
}
```

Benchmark	Test Set
evalPoly-6.00	13
compBal-stdin-6.00	52
compDeriv-6.00	103
hangman2-6.00x	218
prodBySum-6.00	268
oddTuples-6.00	344
hangman1-6.00x	351
evalPoly-6.00x	541
compDeriv-6.00x	918
oddTuples-6.00x	1756
iterPower-6.00x	2875
recurPower-6.00x	2938
iterGCD-6.00x	2988



Average Running Time (in s)



Feedback Generated (Percentage)

Invariant Synthesis for Optimization

Work with Alvin Cheung and Shoaib Kamil



Optimization then and now

Naïve source code



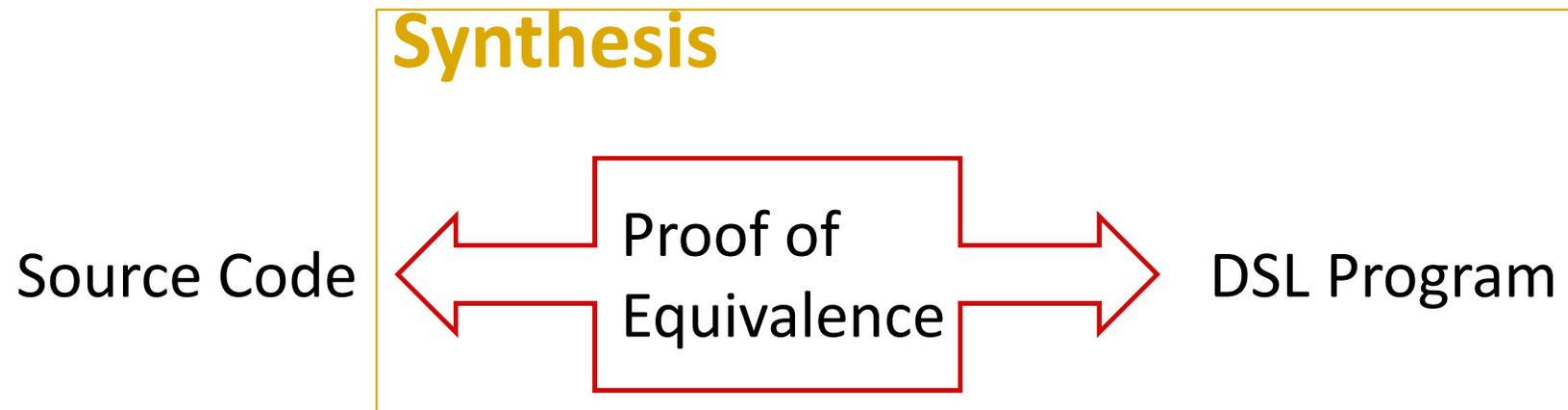
~~Optimal executable~~
Kind-of-OK executable

Domain specific problem description

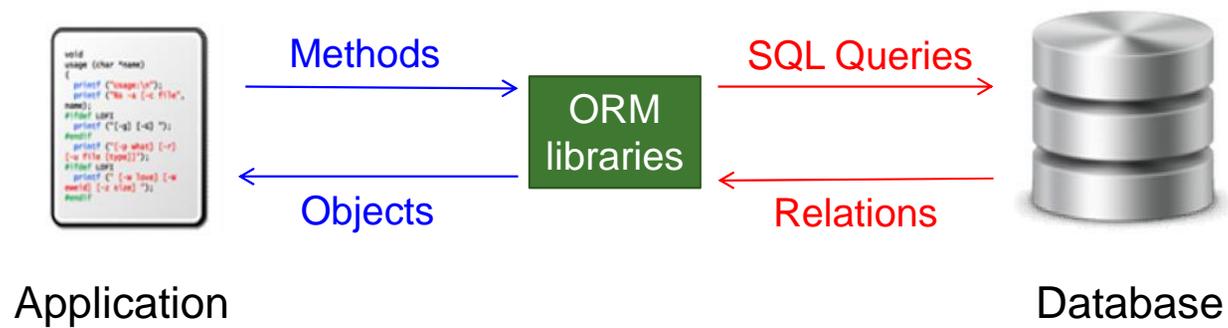


Close to optimal implementation

What about existing
source code?



Java to SQL



Java to SQL

```
List getUsersWithRoles () {  
    List users = User.getAllUsers();  
    List roles = Role.getAllRoles();  
    List results = new ArrayList();  
    for (User u : users) {  
        for (Role r : roles) {  
            if (u.roleId == r.id)  
                results.add(u);  
        }  
    }  
    return results; }  
}
```


convert to

SELECT * FROM user

SELECT * FROM role

```
List getUsersWithRoles () {  
    return executeQuery(  
        "SELECT u FROM user u, role r WHERE u.roleId == r.id  
        ORDER BY u.roleId, r.id"; }  
}
```

Java to SQL

```
List getUsersWithRoles () {  
    List users = User.getAllUsers();  
    List roles = Role.getAllRoles();  
    List results = new ArrayList();  
    for (User u : users) {  
        for (Role r : roles) {  
            if (u.roleId == r.id)  
                results.add(u);  
        }  
    }  
    return results; }  
}
```

outerInvariant(users, roles, u, results, ...)

innerInvariant(users, roles, u, r, results, ...)

results = outputExpr(users, roles)

Verification
conditions

preCondition →

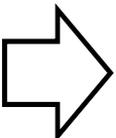
outerInvariant(users/query(...), results/[], ...)

outerInvariant(...) ∧ outer loop terminates →

results = outputExpr(users, roles) ...

Program
+
Unknown Post cond
+
Unknown Invariants

VC
Generator



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C-like language
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$\varphi(c)$

Real-world Evaluation

Wilos (project management application) – 62k LOC

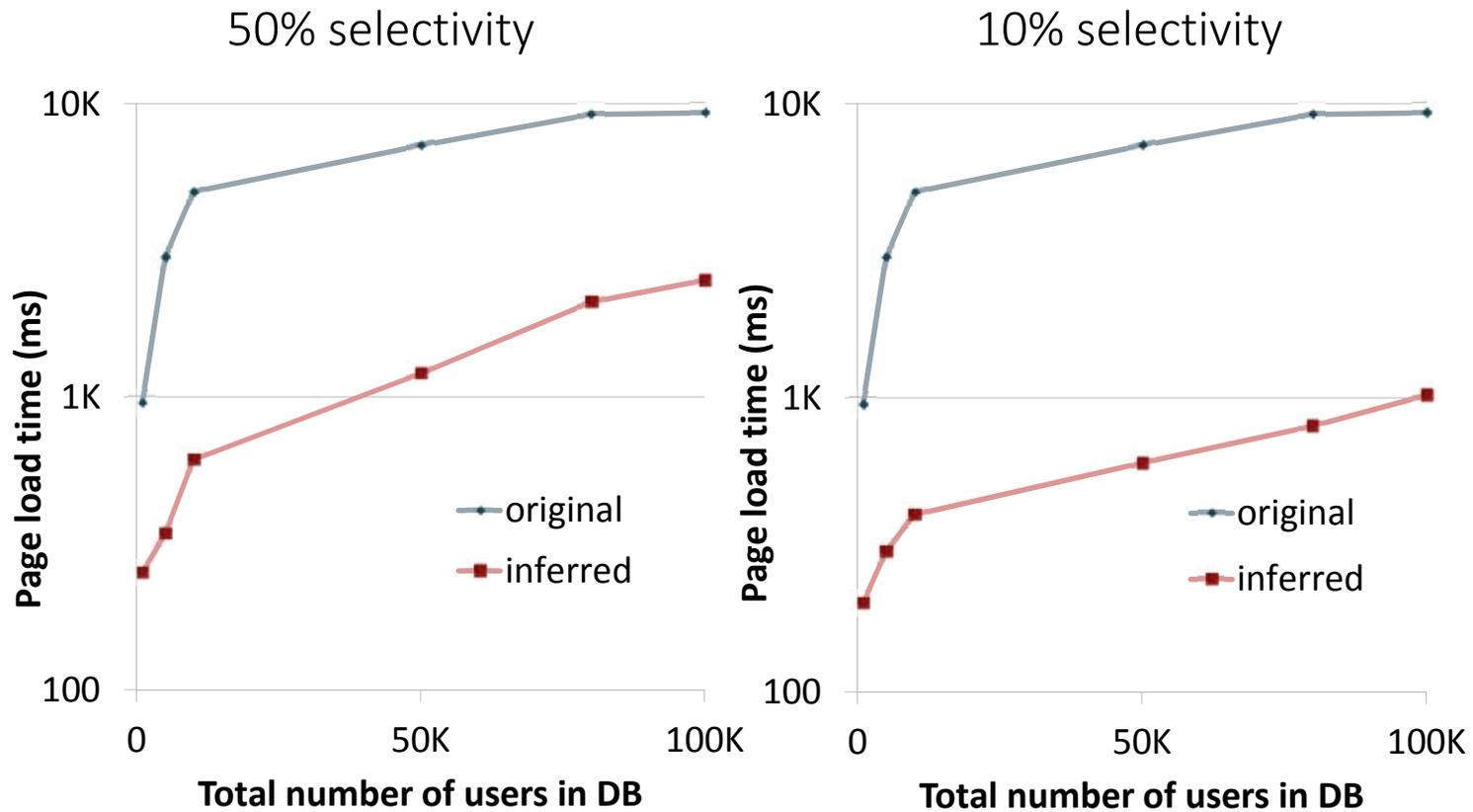
Operation type	# Fragments found	# Fragments converted
Projection	1	1
Selection	13	10
Join	7	7
Aggregation	11	10
Total	33	28

Real-world Evaluation

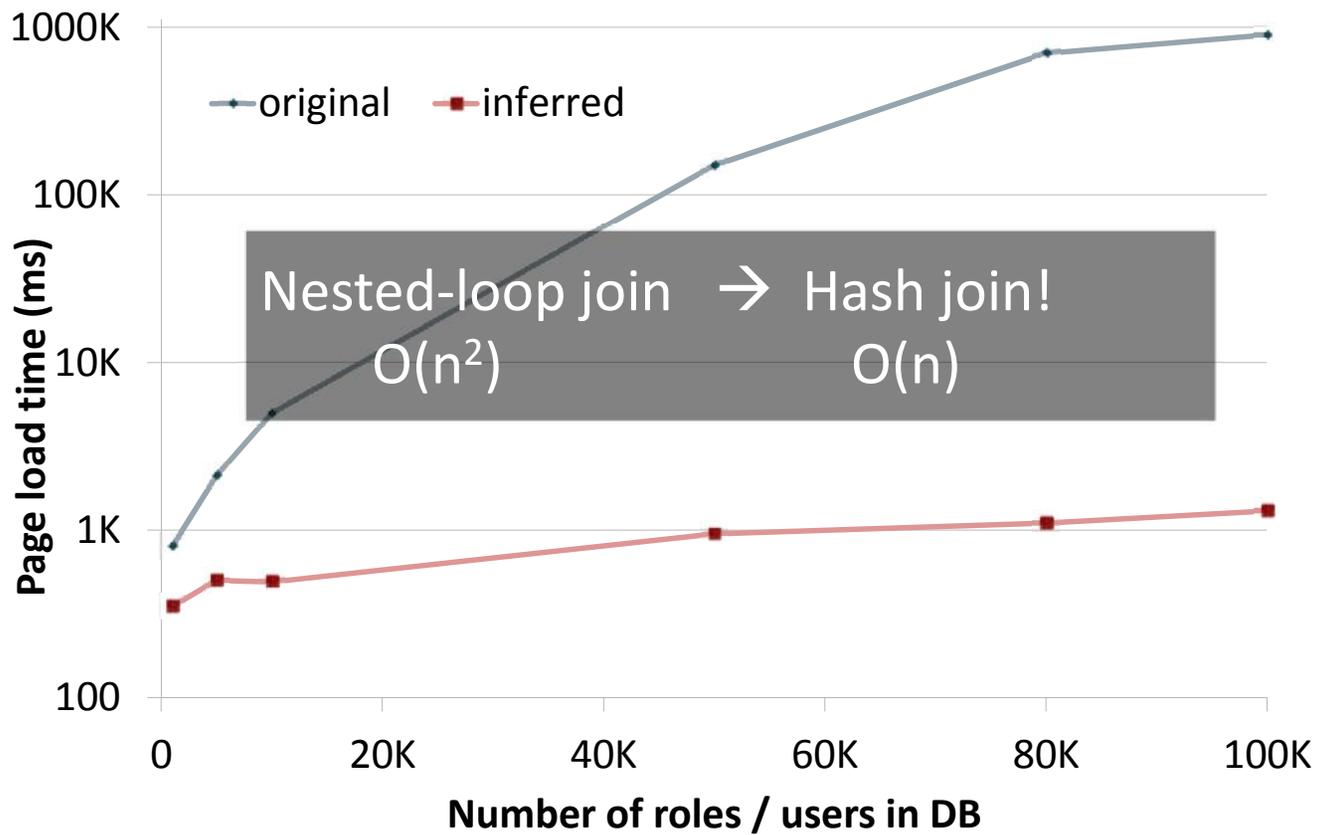
iTracker (bug tracking system) – 61k LOC

Operation type	# Fragments found	# Fragments converted
Projection	3	2
Selection	3	2
Join	1	1
Aggregation	9	7
Total	16	12

Selection Query



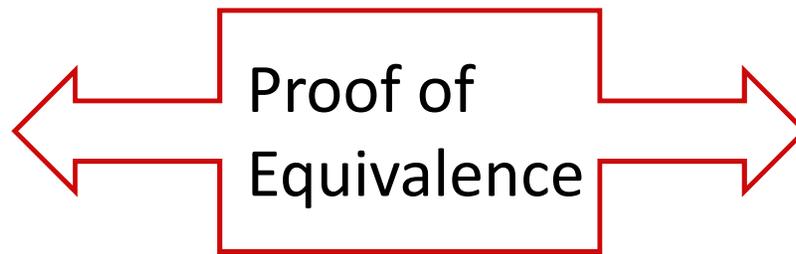
Join Query



What about HPC?

Legacy
Fortran/C++
Code

Synthesis



Stencil DLS
(Halide)

Legacy to Halide

```
for (k=y_min-2;k<=y_max+2;k++) {  
  for (j=x_min-2;j<=x_max+2;j++) {  
    post_vol[((x_max+5)*(k-(y_min-2))+(j)-(x_min-2))]  
      =volume[((x_max+4)*(k-(y_min-2))+(j)-(x_min-2))]  
        + vol_flux_y[((x_max+4)*(k+1 -(y_min-2))+(j)-(x_min-2))]  
        - vol_flux_y[((x_max+4)*(k-(y_min-2))+(j)-(x_min-2))];  
  }  
}
```

$$\forall(j,k) \in Dom \text{ post_vol}[j,k] = \text{volume}[j,k] + \text{vol_flux}[j, k+1] + \text{vol_flux}[j, k]$$

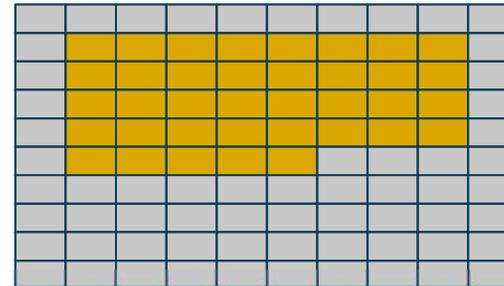
Invariants

$\forall (i, j) \in Dom :$

$$A[i, j] = Expr(\{B_n[expr(i, j), expr(i, j)] \})$$

Invariants

```
for(int i=0; i<n-1; ++i)
  for(int j=0; j<m-1; ++j)
    out[i,j] = sin(in[i,j])
```

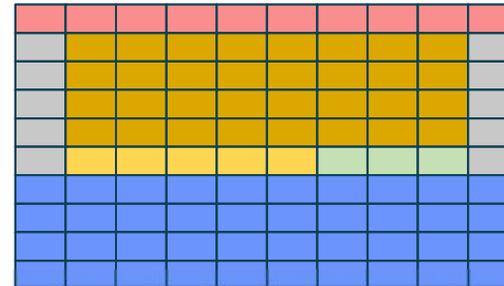


$\forall (i, j) \in Dom :$

$A[i, j] = Expr(\{B_n[expr(i, j), expr(i, j)] \})$

Invariants

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for(int i=0; i<n-1; ++i)
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```



$\forall (i, j) \in Dom :$

$A[i, j] = Expr(\{B_n[expr(i, j), expr(i, j)] \})$

Challenges

- Big invariants
- Complex floating point arithmetic
- Universal Quantifiers

Quantifiers

$$\exists c \forall in P_c(in)$$

```
if(outerInvariant(...) && cond()){  
    assert out == outputExpr(in) ;  
}
```

outerInvariant(...) \wedge outer loop terminates \rightarrow output= outputExpr(in)

The \forall leads to a $\exists \forall \exists$ constraint!

Quantifiers

$$\exists c \forall in P_c(in)$$

Always safe to weaken this condition (more true)

```
if(outerInvariant(...) && cond()){  
    assert out == outputExpr(in) ;  
}
```

$$\bigwedge_{(i,j) \in E} out[i,j] = Expr(\{in[expr(i,j), expr(i,j)] \})$$

Let the synthesizer discover E !

Does it work?

- Can compute invariants for all stencils in Cloverleaf
- Most complex one takes 20 hrs on 15 Cores
- Currently pushing new benchmarks



- There is more to synthesis than “synthesis”
- You can do this too!
 - All the sketch infrastructure is available in open source