

zkay: Specifying and Enforcing Data Privacy in Smart Contracts



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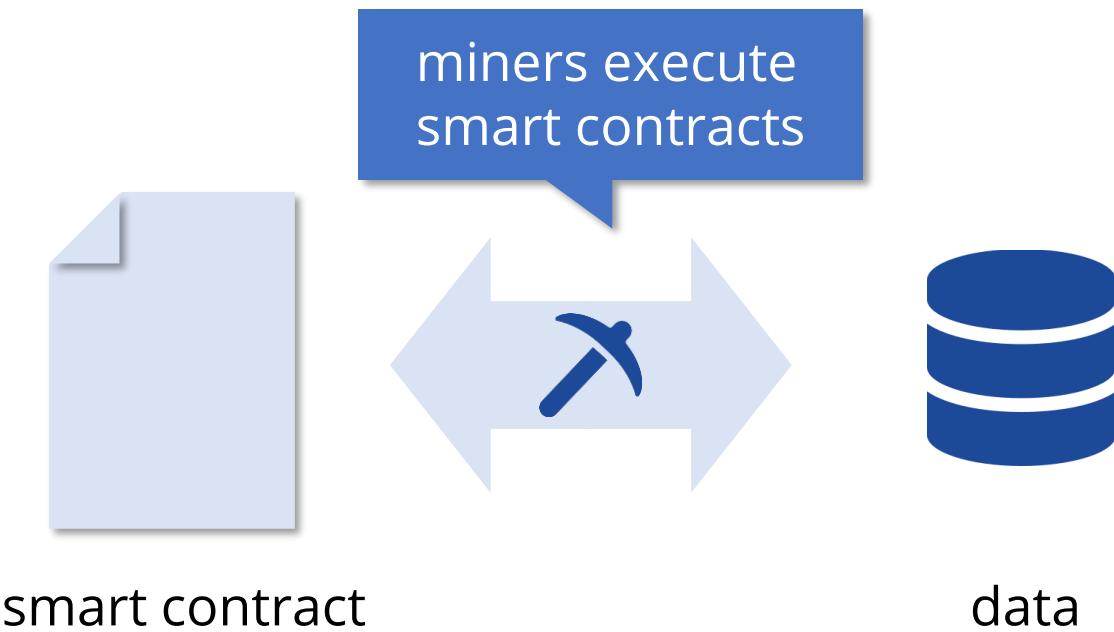


Petar
Tsankov

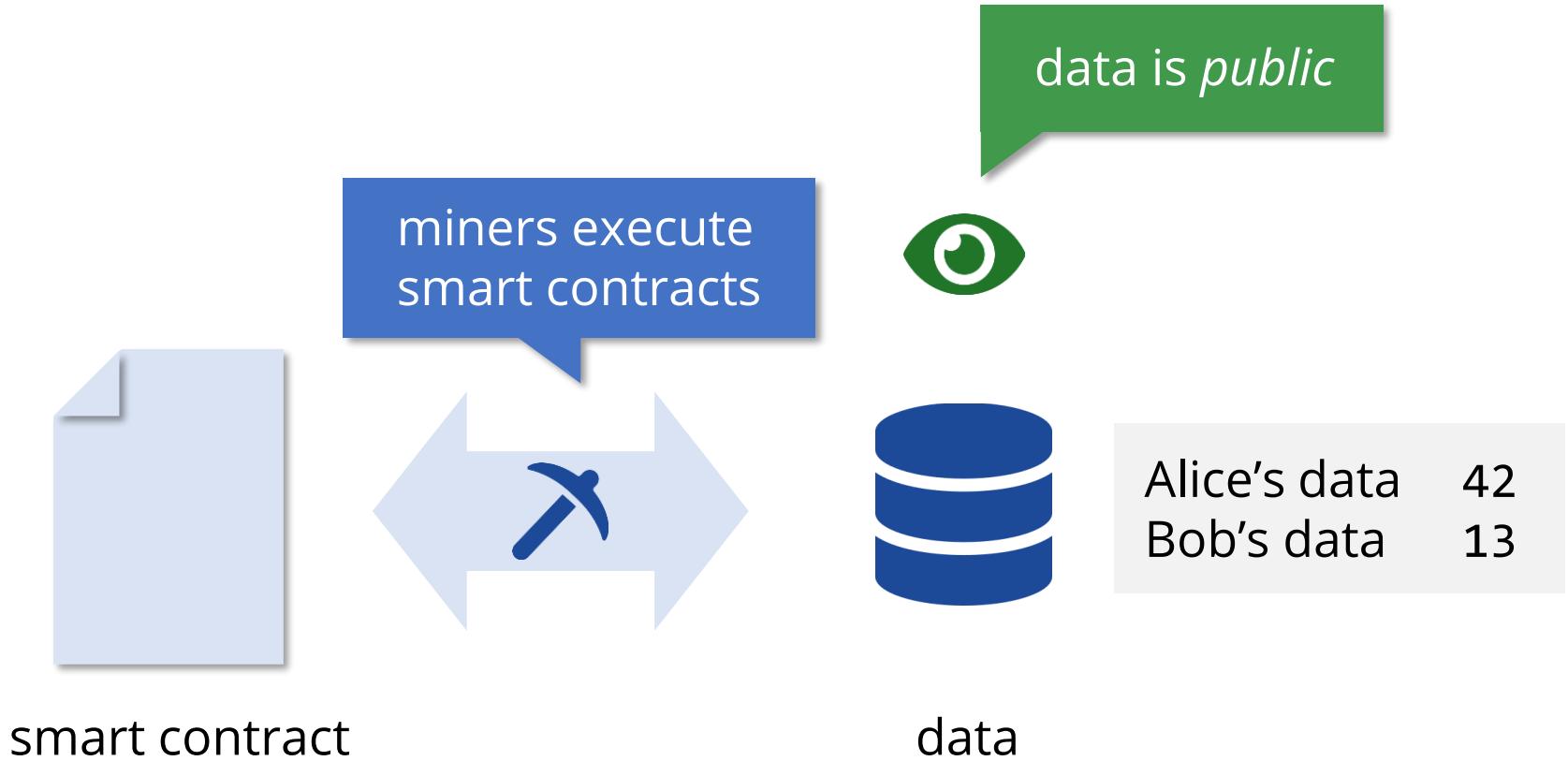


Martin
Vechev

Smart Contracts



Smart Contracts



Privacy?

Medical data and the rise of blockchain

6th July 2018



http://www.pharmatimes.com/web_exclusives/medical_data_and_the_rise_of_blockchain_1243441

Privacy?

Blockchain / Smart Contracts

How Blockchain Could Give Us a Smarter Energy Grid

Energy experts believe that blockchain technology can solve a maze of red tape and data management problems.

by Mike Orcutt

Oct 16, 2017

<https://www.technologyreview.com/s/609077/how-blockchain-could-give-us-a-smarter-energy-grid/>

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4,540 views | Jun 28, 2019, 12:07pm

Paradigm Shift: Biometrics And The Blockchain Will Replace Paper Passports Sooner Than You Think



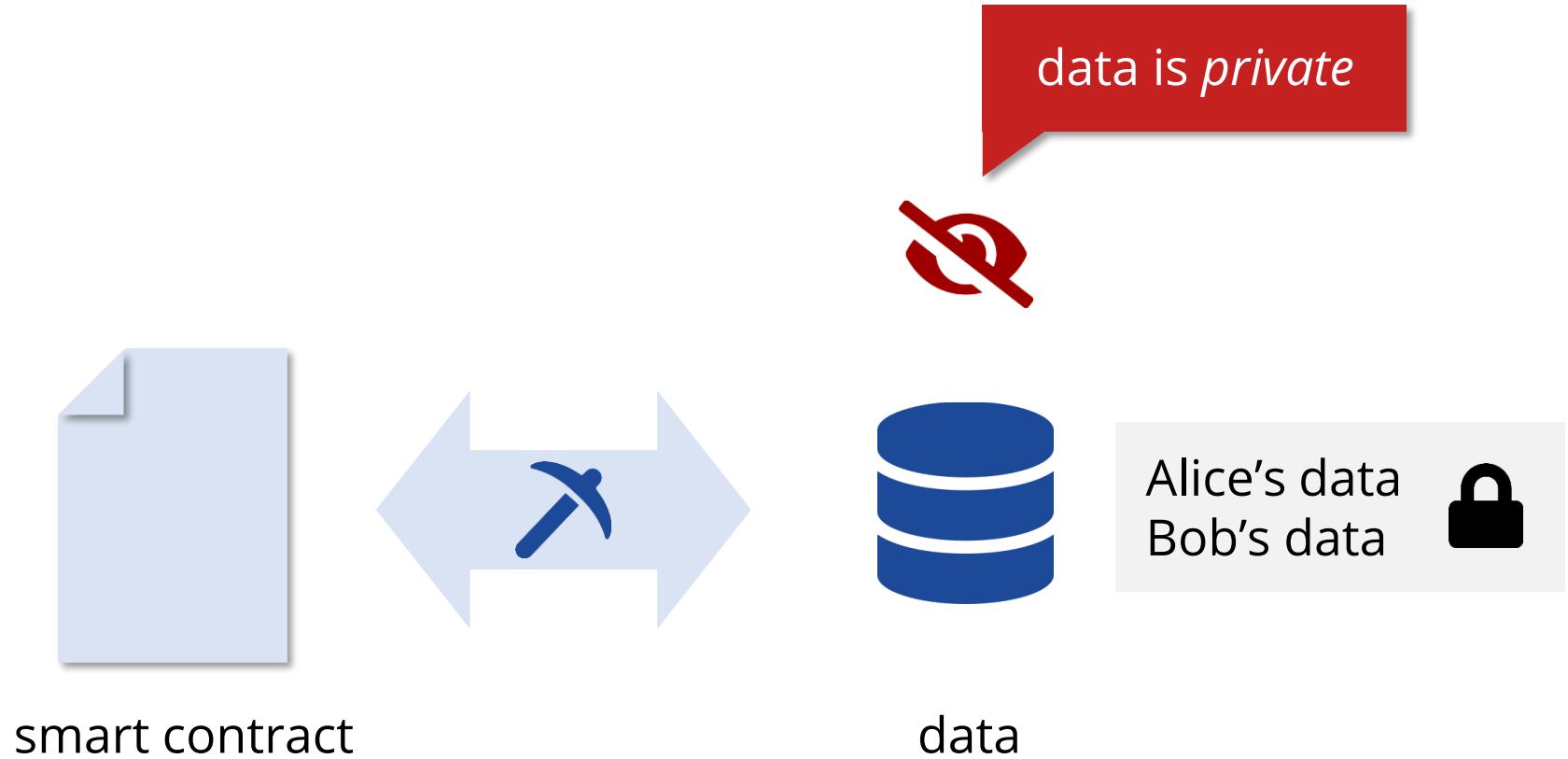
Suzanne Rowan Kelleher Contributor @Travel



Biometrics and blockchain are the keys to the future of traveler identification. GETTY

<https://www.forbes.com/sites/suzannerowankelleher/2019/06/28/paradigm-shift-biometrics-and-the-blockchain-will-replace-paper-passports-sooner-than-you-think/>

Data Privacy for Smart Contracts



Existing Work



- Zerocoin [S&P 13]
- Zerocash [S&P 14]
- Bolt [CCS 17]

Existing Work



Zerocoin [S&P 13]
Zerocash [S&P 14]
Bolt [CCS 17]

only for payments

Existing Work



Zerocoin [S&P 13]
Zerocash [S&P 14]
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only for payments



Hawk [S&P 16]
Arbitrum [Usenix 18]
Ekiden [EuroS&P 19]

Existing Work



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introduce trusted third parties / HW

Existing Work



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introduce trusted third parties / HW

Our work: Leverage *cryptographic primitives* on *existing public* blockchains

Existing Work



Zerocoin [S&P 13]
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Bolt [CCS 17]

only for payments



Hawk [S&P 16]
Arbitrum [Usenix 18]
Ekiden [EuroS&P 19]

introduce trusted third parties / HW

Our work:

Leverage *cryptographic primitives* on *existing public* blockchains

asymmetric encryption



NIZK proofs



Data Privacy with NIZK Proofs

42

Alice's data 0x3025...
Bob's data 0x2543...

Data Privacy with NIZK Proofs

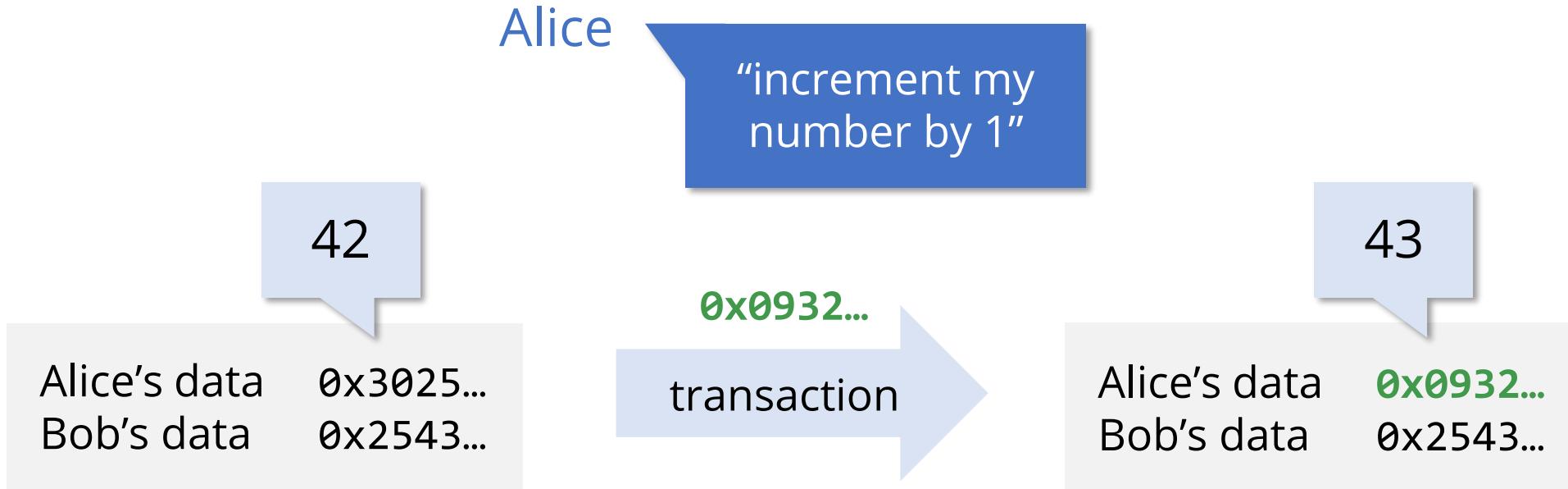
Alice

“increment my
number by 1”

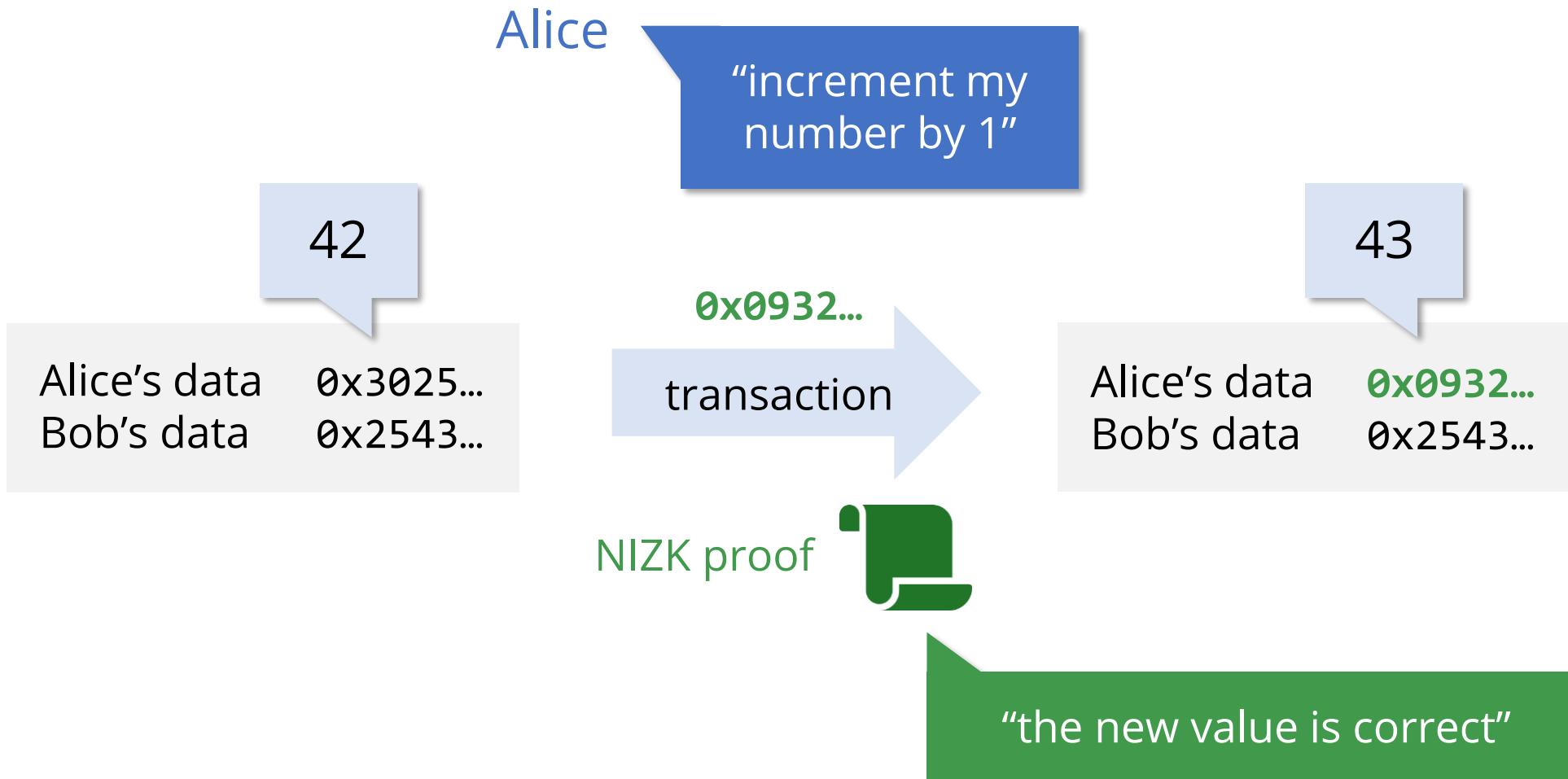
42

Alice's data 0x3025...
Bob's data 0x2543...

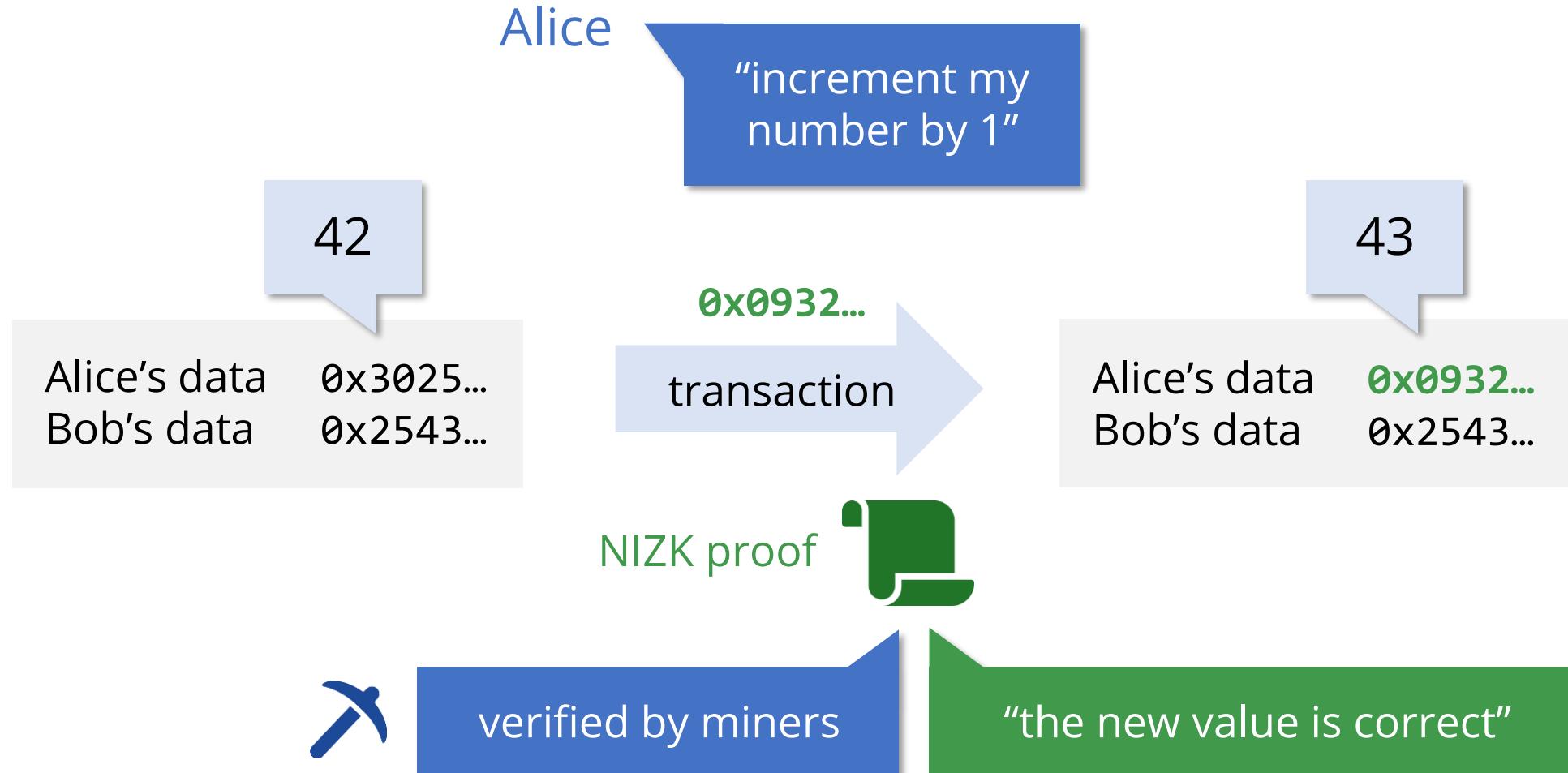
Data Privacy with NIZK Proofs



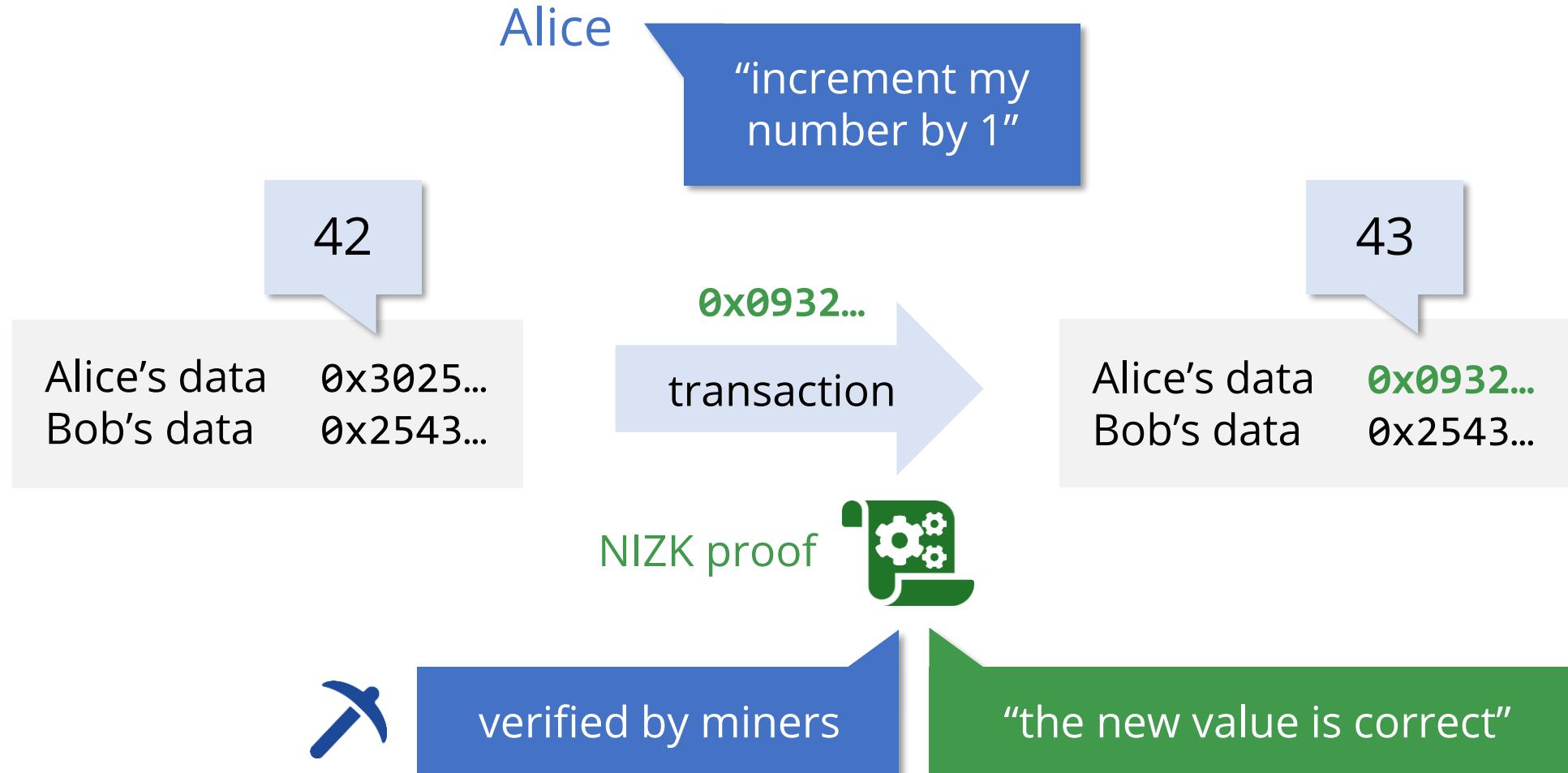
Data Privacy with NIZK Proofs



Data Privacy with NIZK Proofs



Data Privacy with NIZK Proofs



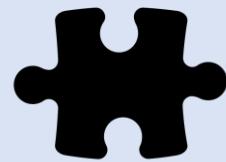
Challenges

Incompleteness of
NIZK Proofs



Challenges

Incompleteness of
NIZK Proofs



e.g., can only incorporate
fixed-size state

Challenges

Incompleteness of
NIZK Proofs



e.g., can only incorporate
fixed-size state

Dynamic arrays?



Entry *lookup*
in contract



Entry *decryption*
in proof

Challenges

Incompleteness of
NIZK Proofs



e.g., can only incorporate
fixed-size state

scattered
logic

Dynamic arrays?



Entry *lookup*
in contract



Entry *decryption*
in proof

Challenges

Incompleteness of
NIZK Proofs



Obfuscated Logic



e.g., can only incorporate
fixed-size state

Dynamic arrays?



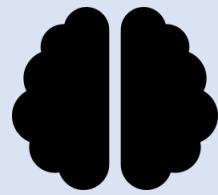
Entry *lookup*
in contract



Entry *decryption*
in proof

Challenges

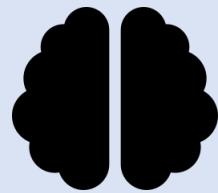
Knowledge
Restrictions



Alice: "increase Bob's value by 1"

Challenges

Knowledge
Restrictions



Alice: "increase Bob's value by 1"

can not generate proof

Challenges

Obfuscated
Information Leaks



Alice: "overwrite Bob's value with my value"

Challenges

Obfuscated
Information Leaks



Alice: "overwrite Bob's value with my value"

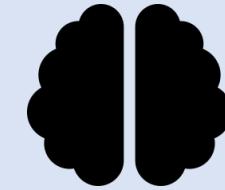
did Alice really want this?

Challenges

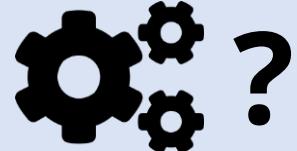
Incompleteness of
NIZK Proofs



Knowledge
Restrictions



Obfuscated Logic



Obfuscated
Information Leaks

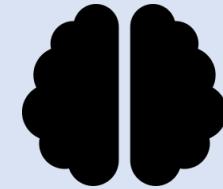


Challenges

Incompleteness of
NIZK Proofs



Knowledge
Restrictions



This work: zkay

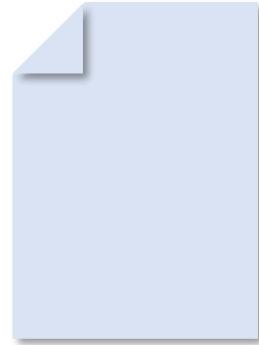
Obfuscated Logic



Obfuscated
Information Leaks

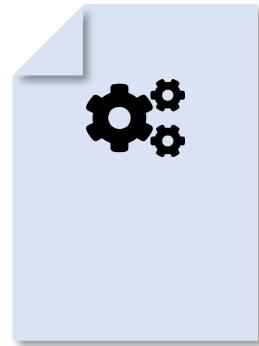


zkay Overview



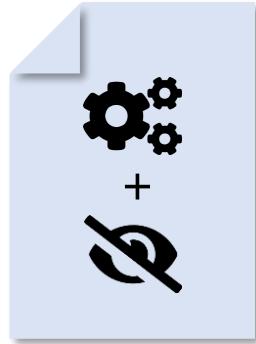
zkay contract

zkay Overview



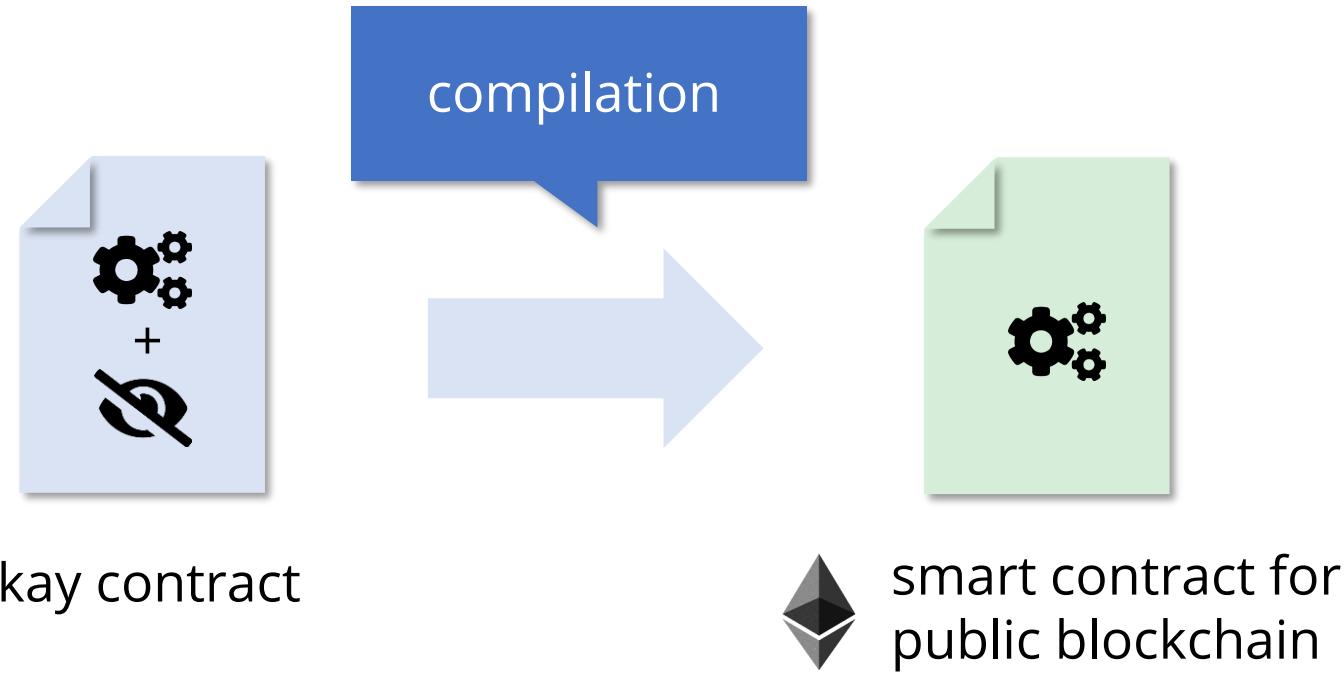
zkay contract

zkay Overview

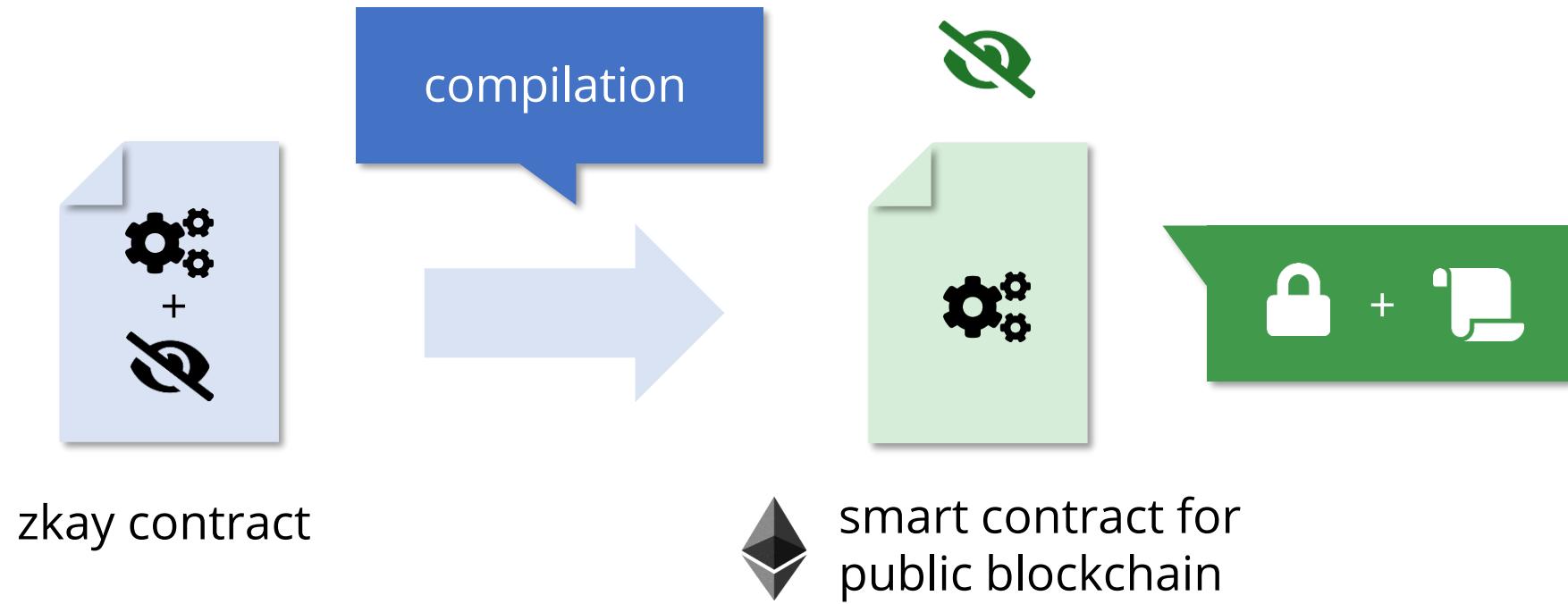


zkay contract

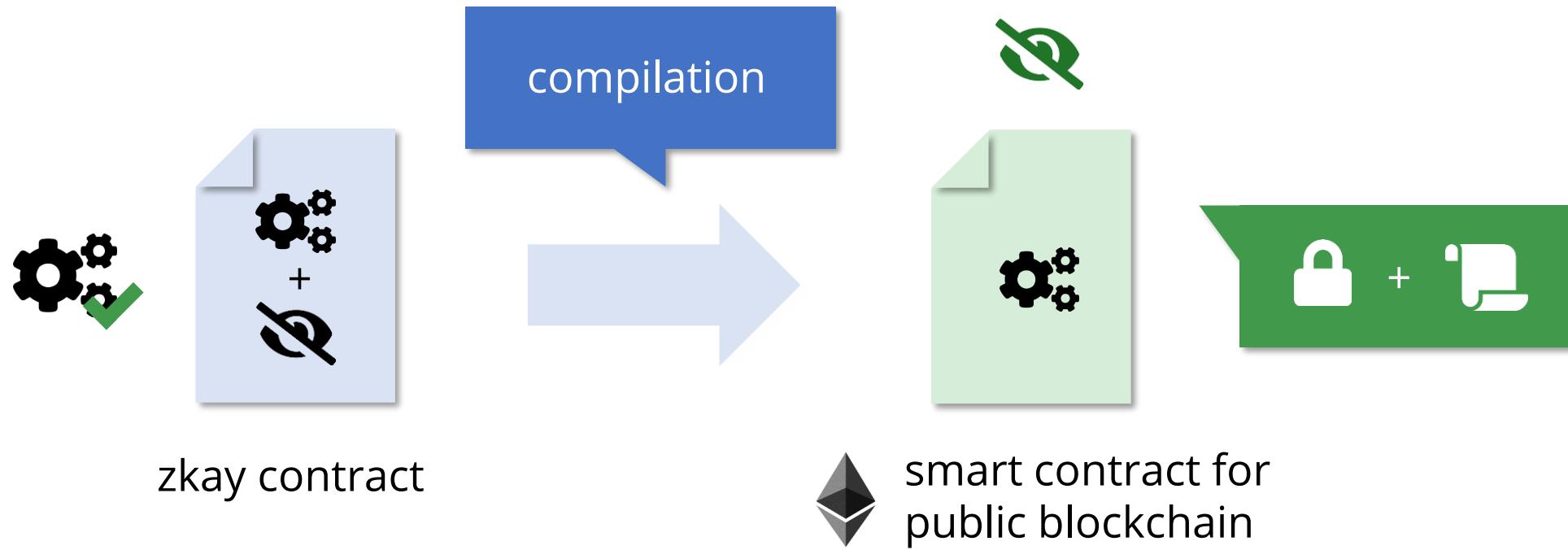
zkay Overview



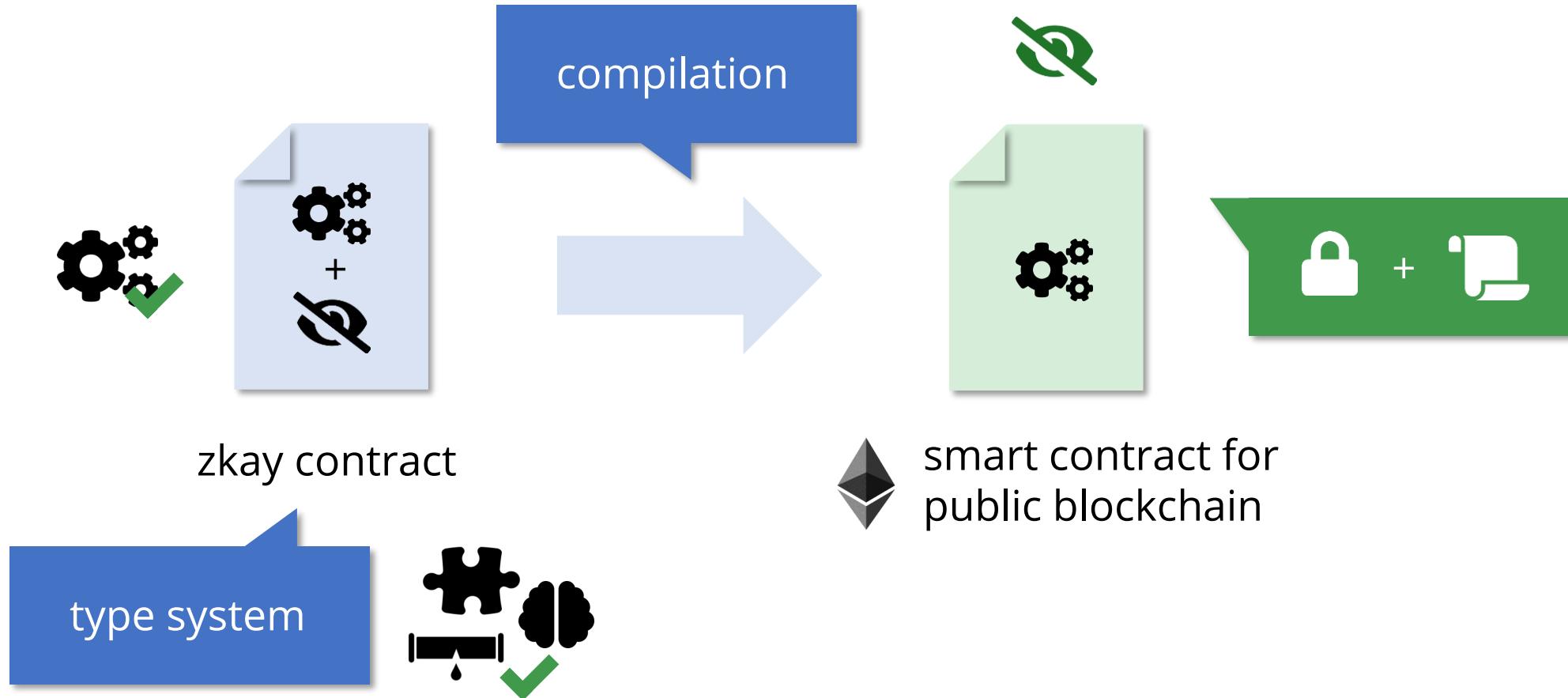
zkay Overview



zkay Overview



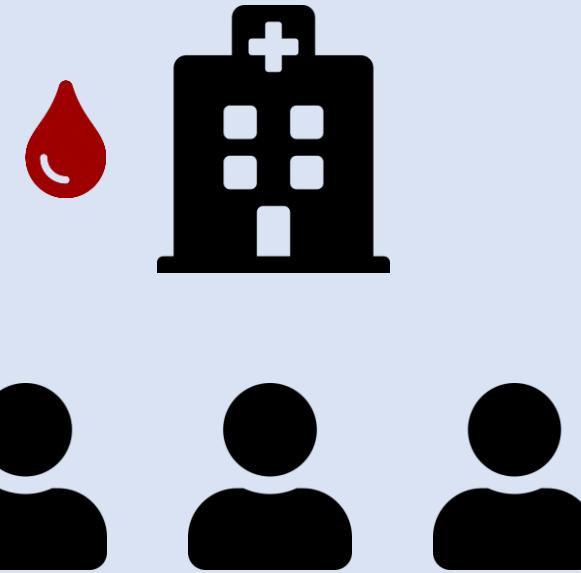
zkay Overview



Example

```
contract MedStats {  
    address hospital;  
    uint count;  
    mapping(address => bool) risk;  
  
    constructor() {  
        hospital = msg.sender; count = 0;  
    }  
  
    function record(address don, bool r) {  
        require(hospital == msg.sender);  
        risk[don] = r;  
        count = count + (r ? 1 : 0);  
    }  
}
```

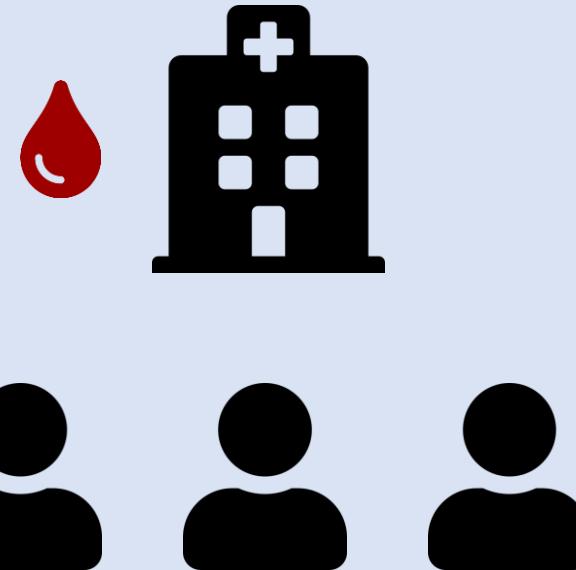
Medical Statistics



Example

```
contract MedStats {  
    address hospital;  
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    mapping(address => bool) risk;  
  
    constructor() {  
        hospital = msg.sender; count = 0;  
    }  
  
    function record(address don, bool r) {  
        require(hospital == msg.sender);  
        risk[don] = r;  
        count = count + (r ? 1 : 0);  
    }  
}
```

Medical Statistics



risk

! !

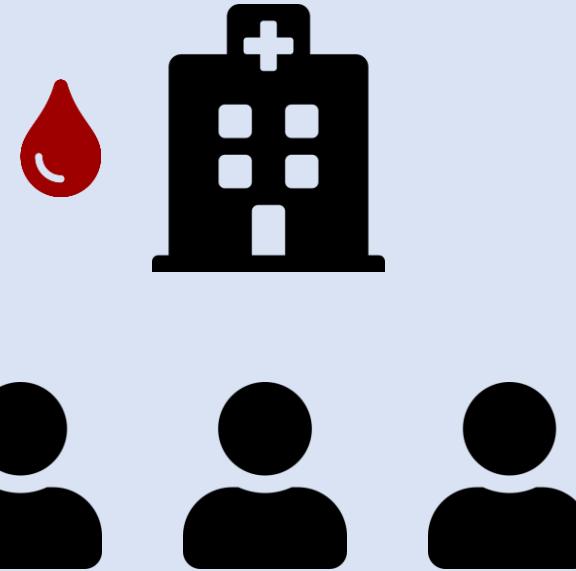
count: 2

Example

Solidity

```
contract MedStats {  
    address hospital;  
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    mapping(address => bool) risk;  
  
    constructor() {  
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    }  
  
    function record(address don, bool r) {  
        require(hospital == msg.sender);  
        risk[don] = r;  
        count = count + (r ? 1 : 0);  
    }  
}
```

Medical Statistics



risk

!

count: 2

Example

Solidity

```
contract MedStats {  
    address hospital;  
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    mapping(address => bool) risk;  
  
    constructor() {  
        hospital = msg.sender; count = 0;  
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    }  
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```

Initialization

Example

Solidity

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contract MedStats {  
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    mapping(address => bool) risk;  
  
    constructor() {  
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    }  
  
    function record(address don, bool r) {  
        require(hospital == msg.sender);  
        risk[don] = r;  
        count = count + (r ? 1 : 0);  
    }  
}
```

Record data of donor

Example

Solidity

```
contract MedStats {  
    address hospital;  
    uint count;  
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Record data of donor

Example

Solidity

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Record data of donor

Example

Solidity

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Record data of donor

Privacy?

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contract MedStats {  
    address hospital;  
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    }  
  
    function record(address don, bool r) {  
        require(hospital == msg.sender);  
        risk[don] = r;  
        count = count + (r ? 1 : 0);  
    }  
}
```

Sensitive information
must remain *private*

zkay Privacy Annotations

Idea:

datatype@owner

zkay Privacy Annotations

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

count is *private to hospital*
hospital is the *owner* of count

zkay Privacy Annotations

```
contract MedStats {  
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    mapping(address!x => bool@x) risk;  
  
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        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

count is *private to hospital*
hospital is the *owner* of count

risk[a] is private to address a

zkay Privacy Annotations

```
contract MedStats {  
    final address hospital;  
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    mapping(address!x => bool@x) risk;  
  
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        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

count is *private to hospital*
hospital is the *owner* of count

risk[a] is private to address a

r is private to the caller (@me)

zkay Privacy Annotations

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
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    constructor() {  
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    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

count is *private* to hospital
hospital is the *owner* of count

risk[a] is private to address a

r is private to the caller (@me)

don is public (@all)

zkay Privacy Annotations

```
contract MedStats {  
    final address hospital;  
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    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

count is *private* to hospital
hospital is the *owner* of count

risk[a] is private to address a

r is private to the caller (@me)

don is public (@all)

reclassify r for don

zkay Privacy Annotations

```
contract MedStats {
    final address hospital;
    uint@hospital count;
    mapping(address!x => bool@x) risk;

    constructor() {
        hospital = me; count = 0;
    }

    function record(address don, bool@me r) {
        require(hospital == me);
        risk[don] = reveal(r, don);
        count = count + (r ? 1 : 0);
    }
}
```

Logic *not obfuscated*



Type System Goals

No unintended information leaks

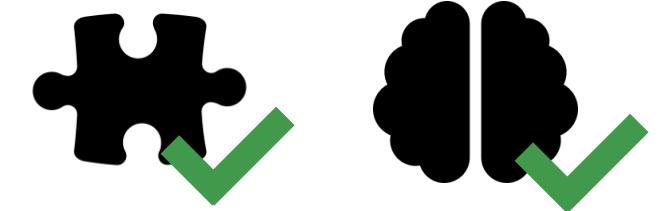


Type System Goals

No unintended information leaks



Realizability using NIZK proofs



Preventing Unintended Leaks

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
        @me          @all  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

Preventing Unintended Leaks

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
    @me  
    @all  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

OK

Allow implicit
classification

Preventing Unintended Leaks

```
contract MedStats {
    final address hospital;
    uint@hospital count;
    mapping(address!x => bool@x) risk;

    constructor() {
        hospital = me; count = 0;
    }

    function record(address don, bool@me r) {
        require(hospital == me);
        risk[don] = r;
        count = count + (r ? 1 : 0);
    }
}
```



The code defines a contract named MedStats. It contains a final variable hospital, a uint variable count, and a mapping risk from address to bool. The constructor initializes hospital to the calling account (me) and count to 0. The record function takes an address don and a bool value r, requiring the calling account to be the same as hospital. It then updates the risk mapping and increments the count by 1 if r is true. Two orange callout boxes point to the 'don' parameter in the record function and the 'me' variable in the require statement.

Preventing Unintended Leaks

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me); disallow  
        risk[don] = r;  
        count = count + (r ? 1 : 0);  
    }  
}
```



Disallow implicit
re-/declassification

Preventing Unintended Leaks

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r  
    }  
}
```

OK

requires **reveal**

Disallow implicit
re-/declassification

Preventing Unintended Leaks

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```



@me

@all

Preventing Unintended Leaks

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
}
```

The code shows three annotations pointing to specific parts of the `record` function:

- A green callout points to the `bool@me r` parameter with the text `@me`.
- A gold callout points to the `address don` parameter with the text `@me`.
- A gold callout points to the `count = count + (r ? 1 : 0);` statement with the text `@all`.

Type of compound
expressions is
conservative

Ensuring Realizability

Alice

```
data[bob] = data[bob] + 1;
```

Ensuring Realizability

Alice

```
data[bob] = data[bob] + 1;
```

@bob

Ensuring Realizability

Alice

```
data[bob] = data[bob] + 1;
```

@bob

can only read @me or @all



Ensuring Realizability

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
    @hospital  
}
```

Can only read *self-owned* or *public* variables

Ensuring Realizability

```
contract MedStats {
    final address hospital;
    uint@hospital count;
    mapping(address!x => bool@x) risk;

    constructor() {
        hospital = me; count = 0;
    }

    function record(address don, bool@me r) {
        require(hospital == me);
        risk[don] = r;
        count = count + (r ? 1 : 0);
    }
}
```

@hospital

Can only read *self-owned* or *public* variables

Ensuring Realizability

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
    can prove hospital == me  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
    @hospital
```

Sometimes, *static analysis* is required

Ensuring Realizability

```
contract MedStats {  
    final address hospital;  
    uint@hospital count;  
    mapping(address!x => bool@x) risk;  
  
    constructor() {  
        hospital = me; count = 0;  
    }  
    can prove hospital == me  
  
    function record(address don, bool@me r) {  
        require(hospital == me);  
        risk[don] = reveal(r, don);  
        count = count + (r ? 1 : 0);  
    }  
    @me OK
```

Sometimes, *static analysis* is required

Compilation

Compilation Example

```
function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
}
```

Compilation Example

```
function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
}
```



```
function record(bin r, bin z, bin p) {  
    require(hospital == me);  
    bin old = count; count = z;  
    verify(p, z, r, old, pk(me));  
}
```

Compilation Example



```
function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
}
```



```
function record(bin r, bin z, bin p) {  
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Compilation Example

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function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
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```



```
function record(bin r, bin z, bin p) {  
    require(hospital == me);  
    bin old = count; count = z;  
    verify(p, z, r, old, pk(me));  
}  
public
```



Compilation Example

```
function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
}
```



```
function record(bin r, bin z, bin p) {  
    require(hospital == me);  
    bin old = count; count = z;  
    verify(p, z, r, old, pk(me));  
}
```



public



private

I know sk such that

$$\text{old} + (\quad r \quad ? 1 : 0)$$

Compilation Example

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function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
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```



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function record(bin r, bin z, bin p) {  
    require(hospital == me);  
    bin old = count; count = z;  
    verify(p, z, r, old, pk(me));  
}
```



public

private

I know sk such that

$$\begin{aligned} \text{old} &+ (r ? 1 : 0) \\ \text{dec}(\text{old}, \text{sk}) + (\text{dec}(r, \text{sk}) ? 1 : 0) \end{aligned}$$



Compilation Example

```
function record(bool@me r) {  
    require(hospital == me);  
    count = count + (r ? 1 : 0);  
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```
function record(bin r, bin z, bin p) {  
    require(hospital == me);  
    bin old = count; count = z;  
    verify(p, z, r, old, pk(me));  
}
```



public



private

I know sk such that

$$\begin{aligned} & \text{old} + (r ? 1 : 0) \\ & \text{dec}(\text{old}, \text{sk}) + (\text{dec}(r, \text{sk}) ? 1 : 0) \\ & z == \text{enc}(\text{dec}(\text{old}, \text{sk}) + (\text{dec}(r, \text{sk}) ? 1 : 0), \text{pk}(me)) \end{aligned}$$

Compilation

<pre> function f(..., \mathcal{F}, bin proof) { P; return e; } function $\bar{f}(\dots, \mathcal{F}, \text{bin proof})$ { T(P) verify ϕ (proof, \mathcal{P}); return $T_e(e)$; } </pre> <p>for $\phi: (\mathcal{P}, \mathcal{S}) \rightarrow \{0, 1\}$</p> <p>(a) Transformation of a zkay function.</p>	$T(P_1; P_2) ::= T(P_1); T(P_2)$ (1) $T(L@\alpha = e@\alpha) ::= T_L(L) = T_e(e)$ (we use $T_L = T_e$) (2) $T(L@\alpha = e@\text{all}) ::= T_L(L) = \text{out}(e, \alpha)$ ($\alpha \neq \text{all}$) (3) $T(\text{require}(e)) ::= \text{require}(T_e(e))$ (4) $T(\text{while } e \{P\}) ::= \text{while } e \{P\}$ (P is fully public) (5) $T(\text{if } e \{P_1\} \text{ else } \{P_2\}) ::= \text{if } T_e(e) \{T(P_1, T_e(e))\} \text{ else } \{T(P_2, T_e(!e))\}$ (6)
$\text{out}(e, \alpha) ::= v_i$ (invariant: $e@\text{me}$ or $e@\text{all}$) $\mathcal{F} \leftarrow \mathcal{F}, v_i$ $\mathcal{P} \leftarrow \mathcal{P}, v_i, \text{pk}(\alpha)$ $\mathcal{S} \leftarrow \mathcal{S}, R_i$ $\phi \leftarrow \phi; v_i == \text{enc}(\text{ } T_\phi(e) \text{ }, R_i, \text{pk}(\alpha))$; <p>(b) Transformation $\text{out}(e, \alpha)$. If $\alpha = \text{all}$, the highlighted part is omitted.</p>	$T_e(c) ::= c$ const c (7) $T_e(\text{id}) ::= \text{id}$ var id (\star) (8) $T_e(L[e]) ::= T_L(L)[T_e(e)]$ mapping entry (9) $T_e((e_1 + e_2)@\text{all}) ::= T_e(e_1) + T_e(e_2)$ native functions (10) $T_e(\text{reveal}(e, \alpha)) ::= \text{out}(e, \alpha)$ (invariant: $e@\text{me}$) (11) $T_e(e@\alpha) ::= \text{out}(e, \alpha)$ (invariant: $e@\text{me}$) (12) <p>(d) Transforming statements using T.</p>
$\text{in}(e, \alpha) ::= \text{dec}_\tau(v_i, \text{sk})$ (invariant: $e@\alpha, \alpha \in \{\text{me, all}\}$) add to $T(P)$: $\tau'@\text{all } v_i = T_e(e);$ for $\tau' ::= \begin{cases} \tau & \alpha = \text{all} \\ \text{bin} & \alpha \neq \text{all} \end{cases}$ $\mathcal{P} \leftarrow \mathcal{P}, v_i$ $\mathcal{S} \leftarrow \mathcal{S}, \text{sk}$ <p>(c) Transformation $\text{in}(e, \alpha)$. If $\alpha = \text{all}$, the highlighted part is omitted.</p>	<p>(e) Transforming expressions in zkay using T_e. For private function arguments id, \star adds an additional correctness constraint to ϕ.</p> $T_e(c) ::= c$ const c (13) $T_e(L@\alpha) ::= \text{in}(L, \alpha)$ (invariant: $\alpha \in \{\text{all, me}\}$) (14) $T_\phi(\text{reveal}(e, \alpha)) ::= T_\phi(e)$ (15) $T_\phi(e_1 + e_2) ::= T_\phi(e_1) + T_\phi(e_2)$ native functions (16) <p>(f) Transforming expressions in the proof circuit using T_ϕ.</p>

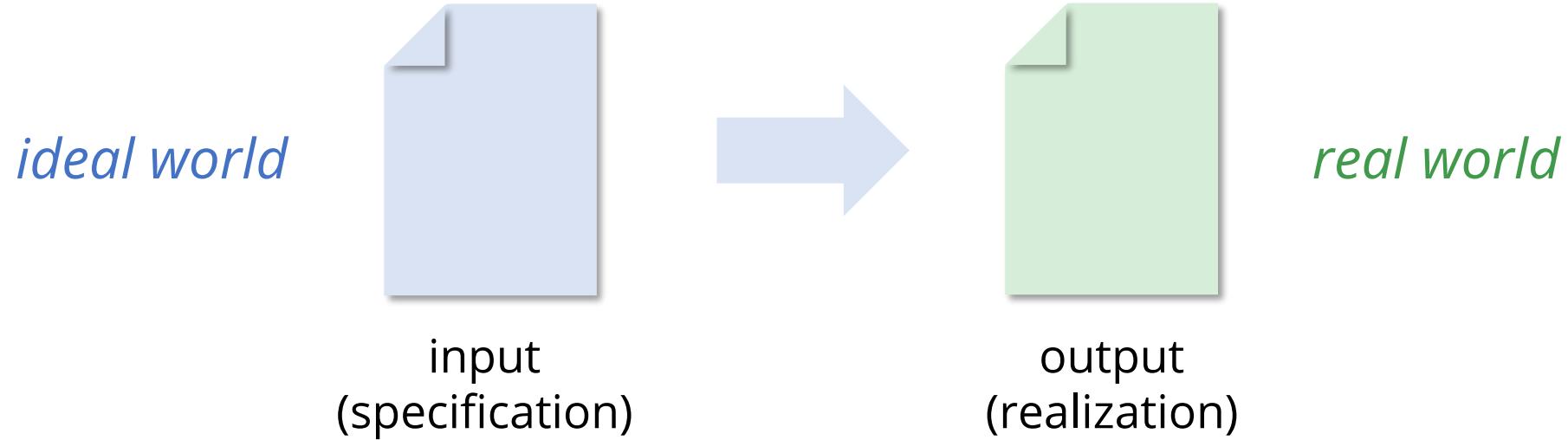
Figure 5: Overview of zkay transformations. We write $e@\alpha$ to indicate that e has privacy type α . The symbol v_i denotes a fresh variable.

see paper for details

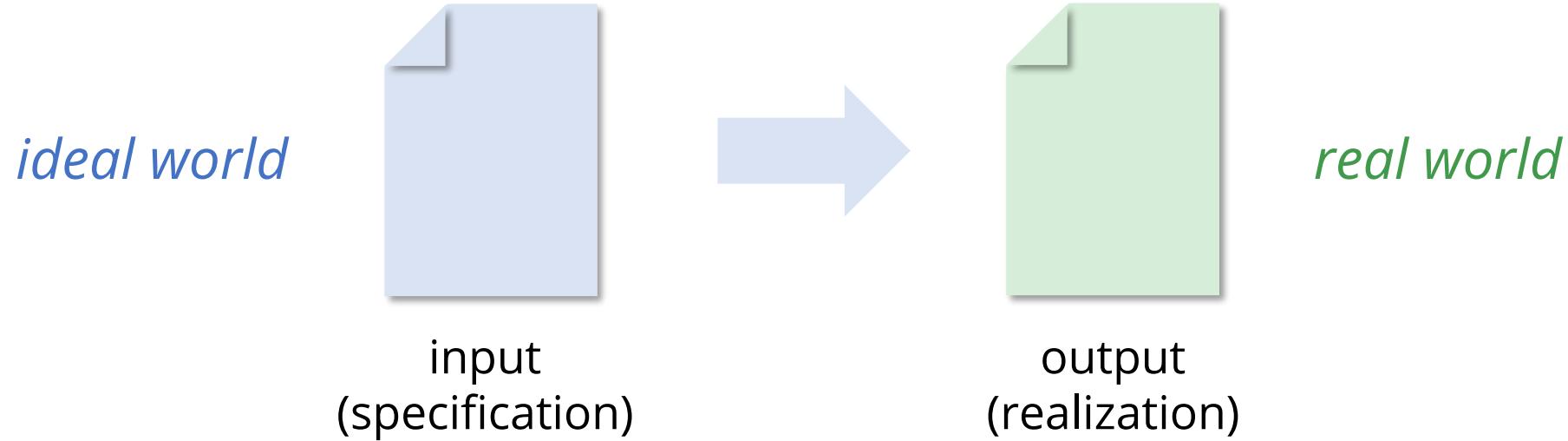
scattered logic

Privacy Guarantees

Data Privacy

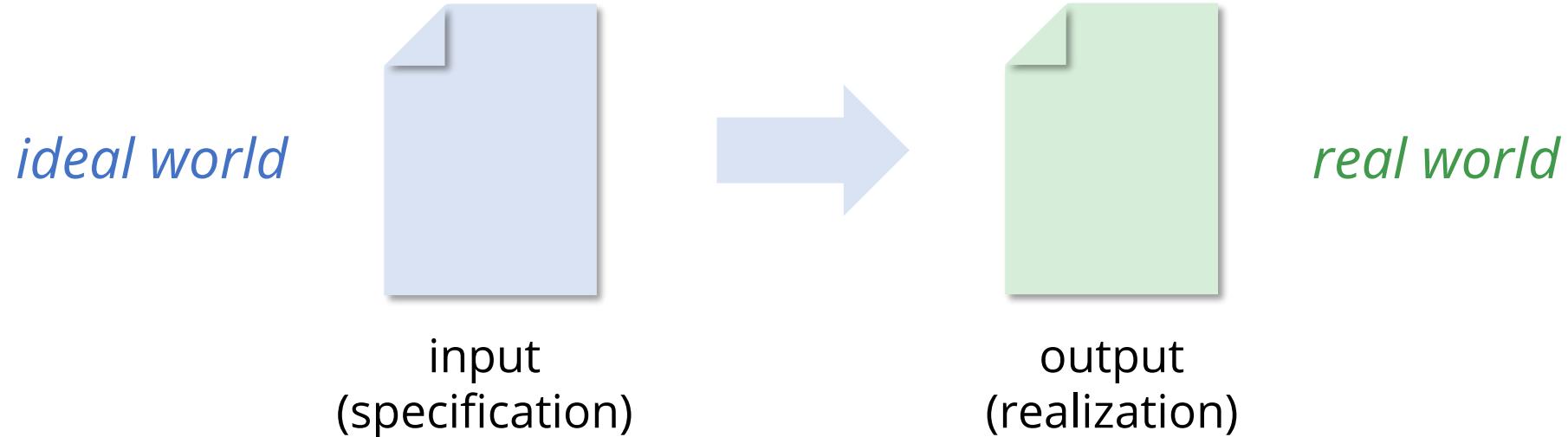


Data Privacy



Everything Eve can
learn here...

Data Privacy

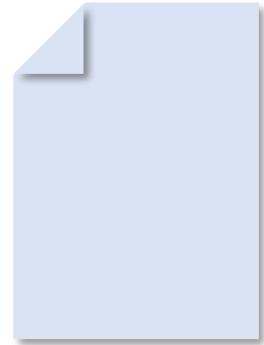


...she can also learn here

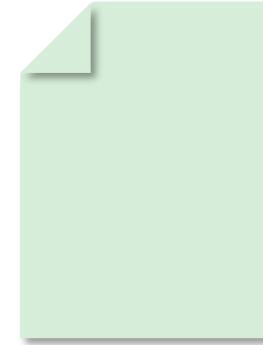
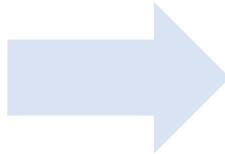
Everything Eve can
learn here...

Data Privacy

ideal world



input
(specification)



output
(realization)

real world

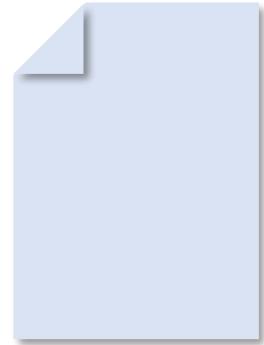
symbolic,
Dolev-Yao style

Everything Eve can
learn here...

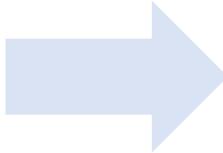
...she can also learn here

Data Privacy

ideal world



input
(specification)



real world



output
(realization)

symbolic,
Dolev-Yao style

Everything Eve can
learn here...

...she can also learn here

according to privacy types

Data Privacy

ideal world



real world

Theorem: Compiled contracts are *private* w.r.t. their input contracts

ic,
'ao style

...she can also learn here

Everything Eve can
learn here...

according to privacy types

Implementation



eth-sri/zkay

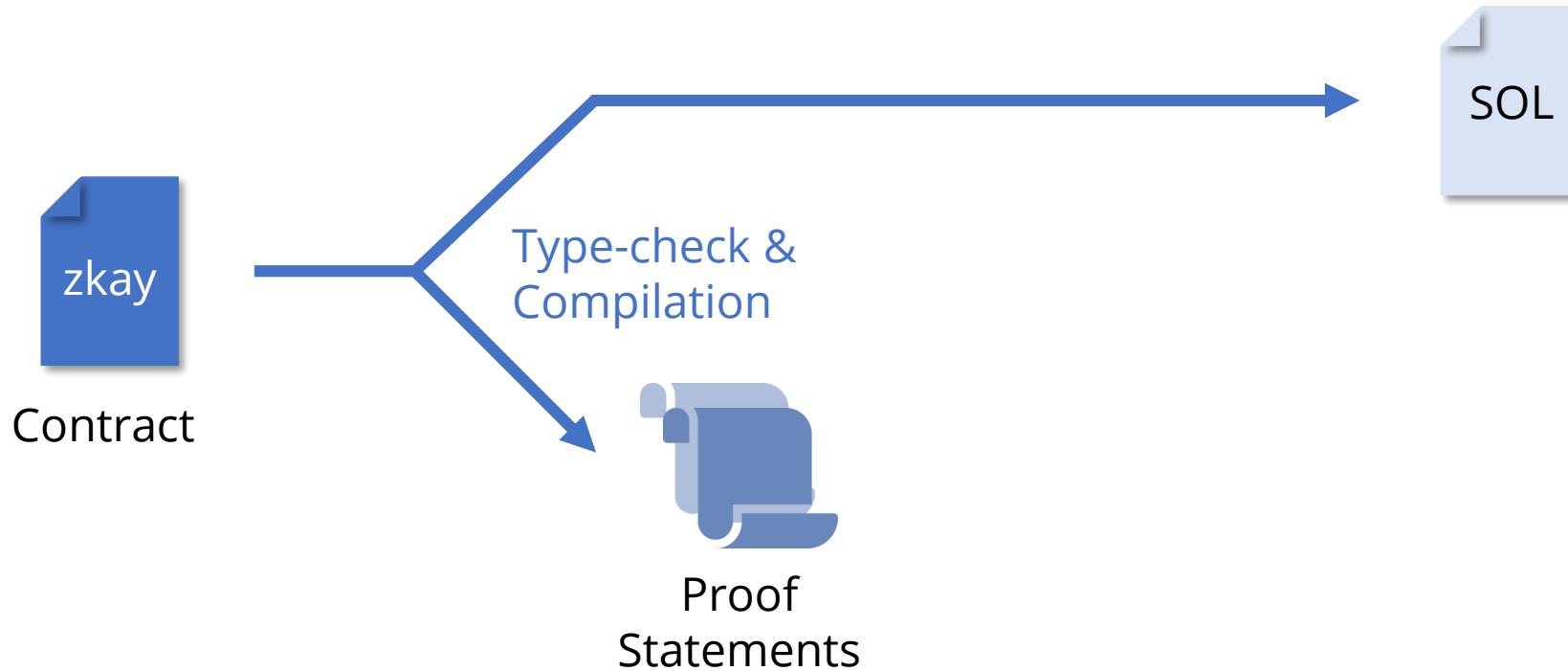


Contract

Implementation



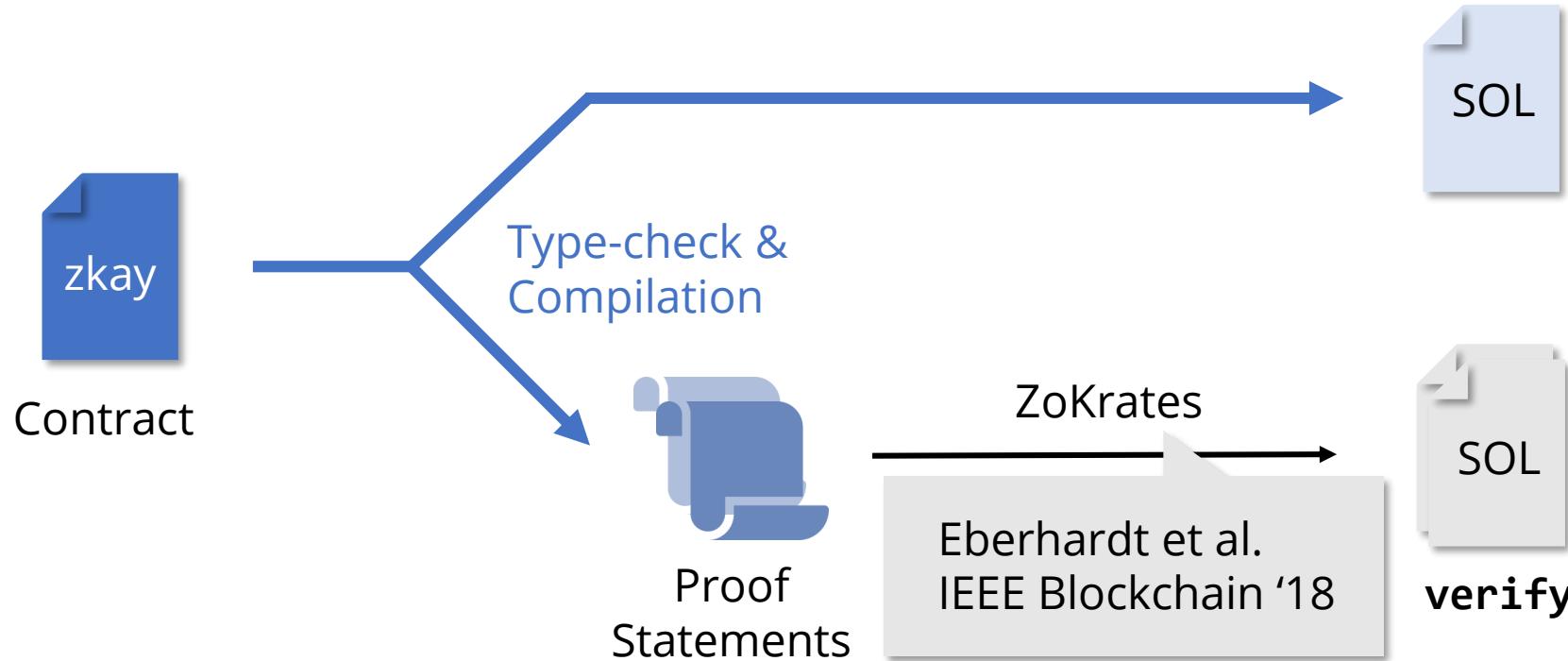
eth-sri/zkay



Implementation



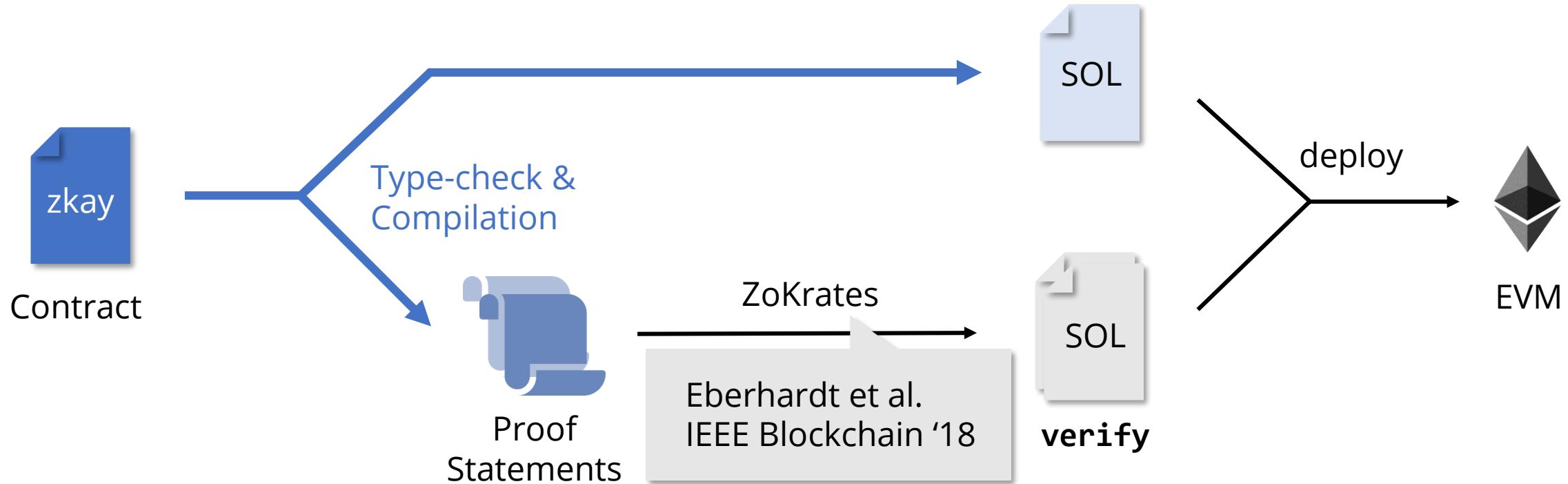
eth-sri/zkay



Implementation



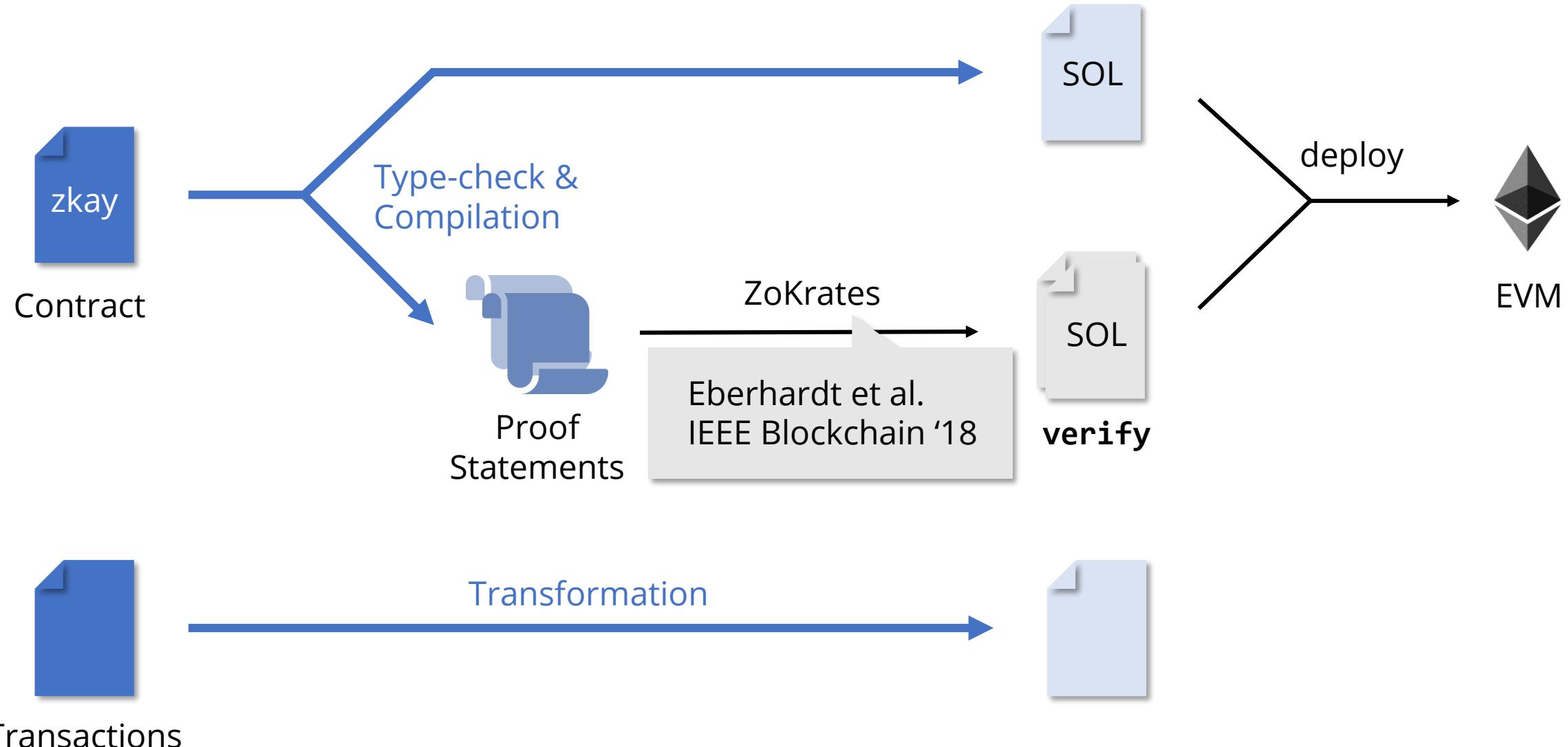
eth-sri/zkay



Implementation



eth-sri/zkay

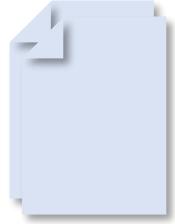


Evaluation

Expressiveness?

Gas costs?

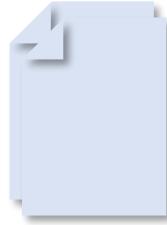
Evaluation



10 Example Contracts

22 – 79 loc

Evaluation



10 Example Contracts

22 – 79 loc

wide range of domains

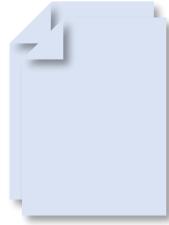
Insurance & Finance

Healthcare

Gambling

...

Evaluation



10 Example Contracts

22 – 79 loc

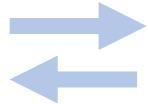
wide range of domains

Insurance & Finance

Healthcare

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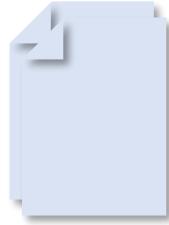
...



Small example scenarios

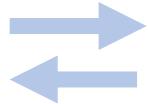
4 – 9 transactions

Evaluation



10 Example Contracts

22 – 79 loc



Small example scenarios

4 – 9 transactions

wide range of domains

Insurance & Finance

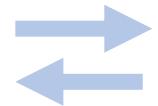
Healthcare

Gambling

...

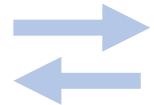
zkay is expressive

On-Chain Costs



Private Transactions

On-Chain Costs

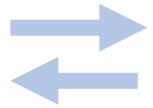


Private Transactions

dominated by constant-cost*
NIZK proof verification

ca. 10^6 Gas per tx

On-Chain Costs



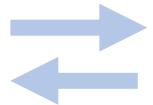
Private Transactions

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ca. 0.50 US\$ (1 ETH = 170 US\$)

On-Chain Costs



Private Transactions

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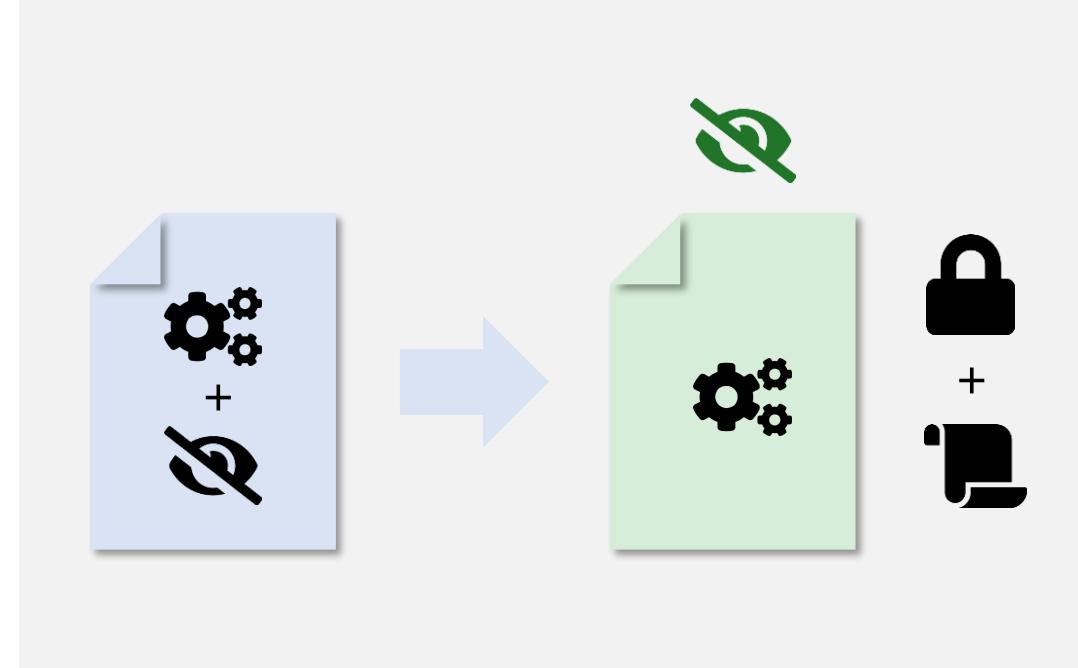
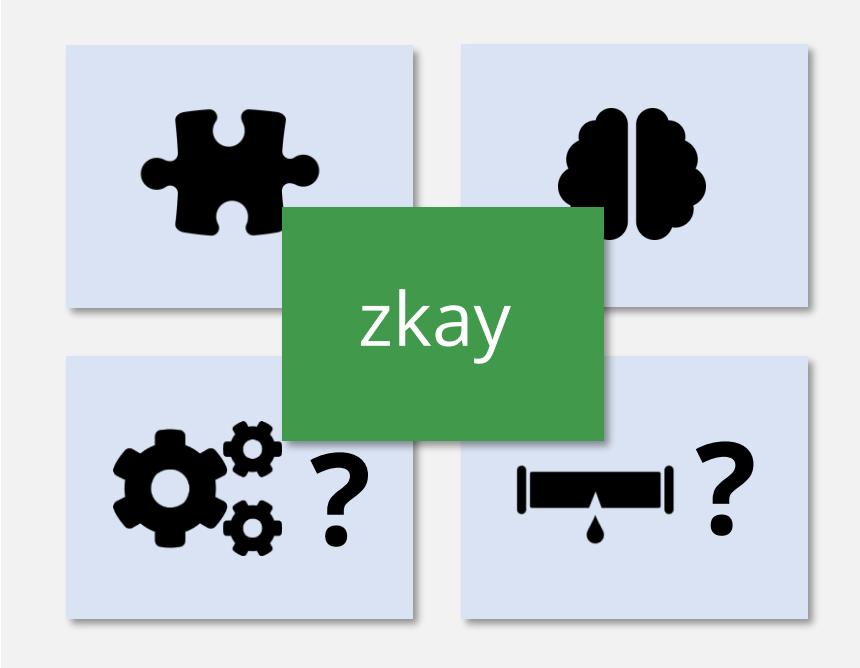
ca. 0.50 US\$ (1 ETH = 170 US\$)

gas costs are moderate

Summary

ETHzürich

SRILAB



datatype@owner



eth-sri/zkay

expressive

moderate costs

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